

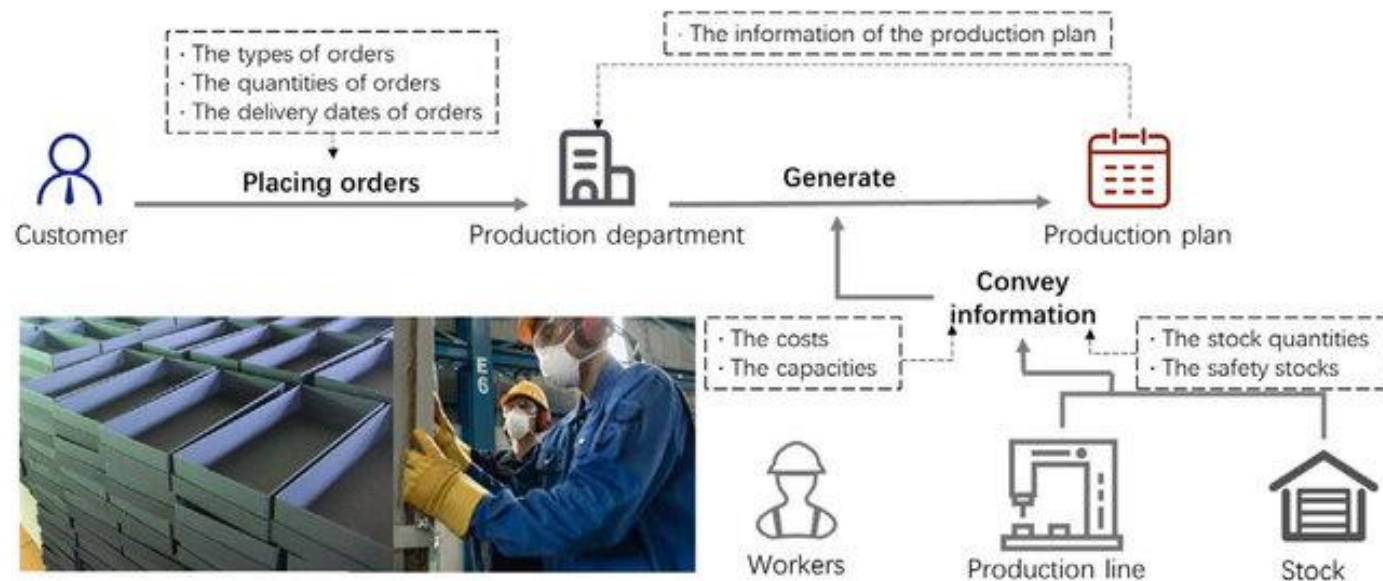
Quantum-Inspired Algorithm for Solving the Production Planning Problem

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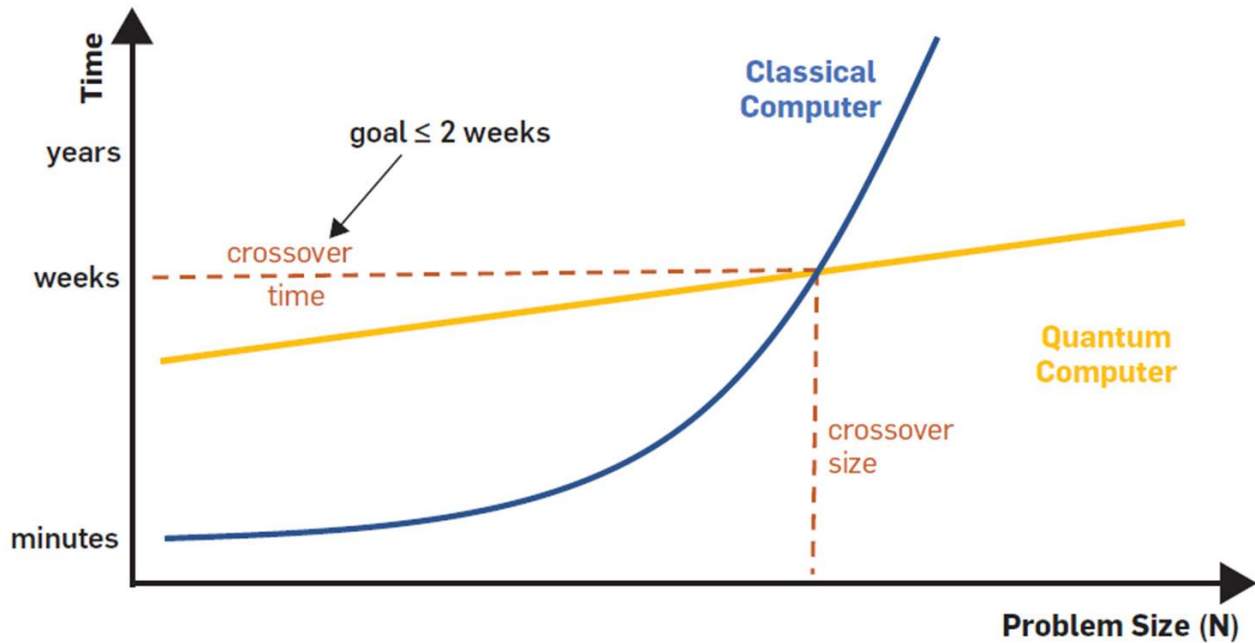
GRID 2025

Production planning

- Input:** customer orders — type, quantity, due-date
- Decision:** production over the planning horizon
- Key limits:** machine capacity, worker shifts & skills, warehouse space / safety stock
- Goal:** minimize total cost of lateness + set-ups + inventory
- Complexity:** reduces to Job/Flow-Shop or RCPSP $\Rightarrow$ **NP-hard**; search space explodes for large plants

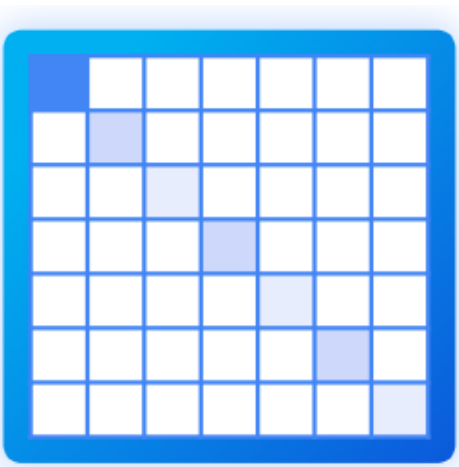


Quantum Computing



- **Explores many solutions at once** via superposition
- **Escapes local minima** better than classical heuristics (quantum tunneling)
- **Good fit for combinatorial problems** like scheduling, routing, packing
- **Scales better** for large instances where classical solvers slow down rapidly
- **Hybrid solvers** already show speed gains on mid-size problems (thousands of variables)

QUBO



Data + constraints + objectives



Decision vars

- **QUBO** = Quadratic Unconstrained Binary Optimization
 $\mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x}$ to **Min**, where $\mathbf{x} \in \{0,1\}^n$
- **All constraints are encoded as penalties** in the objective (e.g., capacity limits, precedence, timing)
- **Universal format** for quantum annealing and hybrid solvers
- Used in scheduling, routing, packing, allocation, etc.

MILP

$$\begin{aligned} \text{Min } & \mathbf{c}^T \mathbf{x} \\ \text{s.t. } & \mathbf{A} \mathbf{x} \leq \mathbf{b} \end{aligned}$$

QUBO

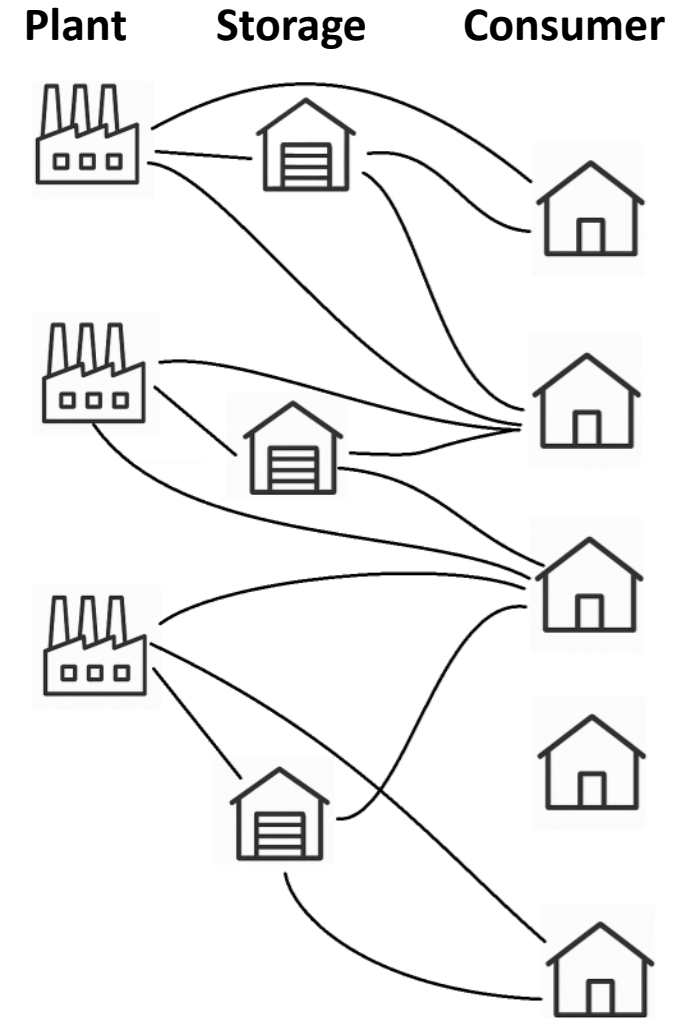
$$\begin{aligned} \rightarrow & \text{Min } \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x} \\ \rightarrow & + \lambda \|\mathbf{A} \mathbf{x} - \mathbf{b}\|^2 \text{ as penalty} \end{aligned}$$

Quantum Algorithms in Production Scheduling

- **Aggoune & Deleplanque (2023)**
Solved job-shop scheduling with 4 jobs $\times$ 4 machines using QUBO on D-Wave.
- **Pérez Armas et al. (2024)**
Solved RCPSP instances with up to 30 activities and 2 resources.
- **Pakhomchik et al. (2023)**
Applied hybrid QUBO method to real automotive workflows with $\sim$ 100 operations.
- **Toma et al. (2023)**
Encoded flexible job-shop problems with up to 40 operations and multiple machines.
- **Zhang et al. (2022)**
Solved nurse scheduling instances with $\sim$ 60 shifts using QUBO + local search on Fujitsu DA.

Production problem

- **Objective:** Develop an optimal production and distribution plan to minimize total costs.
- **Plants:** Multiple with varying production capacities and costs, some existing, others new or upgradable.
- **Consumers:** Several consumers with specific annual demand for products.
- **Storage Constraints:** Products have limited storage times, impacting inventory management.
- **Planning Horizon:** Multi-year period (10+ years).
- **Constraints:**
 - Meet consumer demand annually.
 - Respect plant production capacities and compatibility with consumers.
 - Adhere to storage duration limits and delivery schedules.
 - Account for construction/upgrade costs for new or existing plants.



Plant Capacity Constraint

Ensure production at plant z in year t does not exceed capacity c_z .

MILP Format

- $y_{zt} \leq c_z$, for all z, t
- Vector form: $y_j \leq c_{zj}$, where $j = Z(t - 1) + z$

QUBO Format

- Binarize: $y_{zt} = \sum 2^k y_{jk}$, $y_{jk} \in \{0,1\}$, $k = 0$ to $\lfloor \log_2 c_{zj} \rfloor$
- Penalty: $\rho \sum (\sum 2^k y_{jk} - c_{zj} + \varepsilon_j)^2$, $\varepsilon_j = \text{slack}$

QUBO: $\min y^T Q y + c^T y$, Q encodes penalty

Supply Plan Constraint

Supply from plant z to consumer r in year t must not exceed demand d_{rt} .

MILP Format

- $x_{rzt} \leq d_{rt}$, for all r, z, t
- Vector form: $x_i \leq d_{r_i t_i}$, where $i = RZ(t - 1) + R(z - 1) + r$

QUBO Format

- Binarize: $x_{rzt} = \sum 2^k x_{ik}, x_{ik} \in \{0,1\}, k = 0 \text{ to } \lfloor \log_2 d_{r_i t_i} \rfloor$
- Penalty: $\rho \sum (\sum 2^k x_{ik} - d_{r_i t_i} + \varepsilon_i)^2, \varepsilon_i = \text{slack}$

QUBO: $\min x^T Q x + c^T x, Q \text{ encodes penalty}$

Compatibility Constraint

No supply from plant z to consumer r if incompatible ($p_{rz} = 0$).

MILP Format

- $x_{rzt} = 0$ if $p_{rz} = 0$, for all t
- Vector form: $x_i = 0$ if $p_{r_i z_i} = 0$, $i = RZ(t_i - 1) + R(z_i - 1) + r_i$

QUBO Format

- Binarize: $x_{rzt} = \sum 2^k x_{ik}, x_{ik} \in \{0,1\}$
- Penalty: $\rho \sum_{i: p_{r_i z_i} = 0} \left(\sum 2^k x_{ik} \right)^2$

QUBO: $\min x^T Q x$, Q penalizes non-zero x_{ik} when $p_{r_i z_i} = 0$

Delivery Timing Constraint

No supply from plant z to consumer r during restricted years after start (t in $[g_r, g_r + h_{rz} - 1]$).

MILP Format

- $x_{rzt} = 0$ if $t \in [g_r, g_r + h_{rz} - 1]$, for all r, z
- Vector form: $x_i = 0$ if $t_i \in [g_{r_i}, g_{r_i} + h_{r_i z_i} - 1]$

QUBO Format

- Binarize: $x_{rzt} = \sum 2^k x_{ik}$, $x_{ik} \in \{0,1\}$
- Penalty: $\rho \sum_{i: t_i \in [g_{r_i}, g_{r_i} + h_{r_i z_i} - 1]} \left(\sum 2^k x_{ik} \right)^2$

QUBO: $\min x^T Q x$, Q penalizes non-zero x_{ik} in restricted periods

Production Balance Constraint

Total supply from plant z across all consumers and years equals its total production.

MILP Format

- $\sum_{r,t} x_{rzt} = \sum_t y_{zt}, \text{ for all } z$
- Vector form: $A_1 w = b_1, A_1 = (I_{1 \times T} \otimes E_Z \otimes I_{1 \times R}, -I_{1 \times T} \otimes E_Z), b_1 = 0$

QUBO Format

- Binarize: $x_{rzt} = \sum 2^k x_{ik}, y_{zt} = \sum 2^k y_{jk}$
- Penalty: $\rho_1 \sum_z \left(\sum_i A_{1i} \sum 2^k x_{ik} - \sum_j A_{1j} \sum 2^k y_{jk} \right)^2$

QUBO: $\min w^T Q w, Q = \rho_1 A_1^T A_1, w = (x, y)$

Demand Satisfaction Constraint

Total supply to consumer r in year t meets demand d_{rt} .

MILP Format

- $\sum_z x_{rzt} = d_{rt}, \text{ for all } r, t$
- Vector form: $A_2 w = b_2, A_2 = (E_T \otimes I_{1 \times Z} \otimes E_R, 0_{RT \times ZT}), b_2 = d$

QUBO Format

- Binarize: $x_{rzt} = \sum 2^k x_{ik}$
- Penalty: $\rho_1 \sum_{r,t} (\sum_i A_{2i} \sum 2^k x_{ik} - d_{rt})^2$

QUBO: $\min x^T Q x, Q = \rho_1 A_2^T A_2$

Storage Duration Constraint

Production at plant z up to year t must not exceed supply within storage limit s_z .

MILP Format

- $\sum_{s=1}^t y_{zs} \leq \sum_r \sum_{s=1}^{t+s_z} x_{rzs}, \quad t \leq T - s_z - 1, \text{ for all } z$
- Vector form: $C_2^z w \leq 0, C_2^z = (-M^z \otimes e^z \otimes I_{1 \times R}, M(1:T - s_z - 1, :) \otimes e^z)$

QUBO Format

- Binarize: $x_{rzs} = \sum 2^k x_{ik}, y_{zs} = \sum 2^k y_{jk}$
- Penalty: $\rho_2 \sum_{z,t} \left(\sum_j C_{2j}^z \sum 2^k y_{jk} - \sum_i C_{2i}^z \sum 2^k x_{ik} + \varepsilon_{zt} \right)^2$

QUBO: $\min w^T Q w, Q = \rho_2 (C_2^z)^T C_2^z$

Real problem

- 5 consumers
- 3 plants
- Storage times 1, 3, 5 years

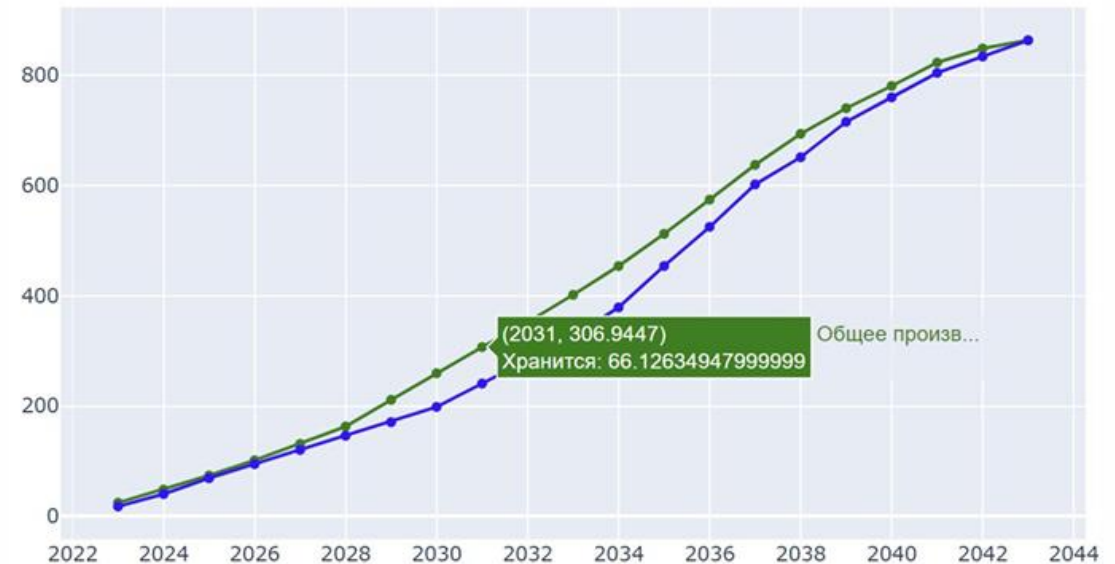
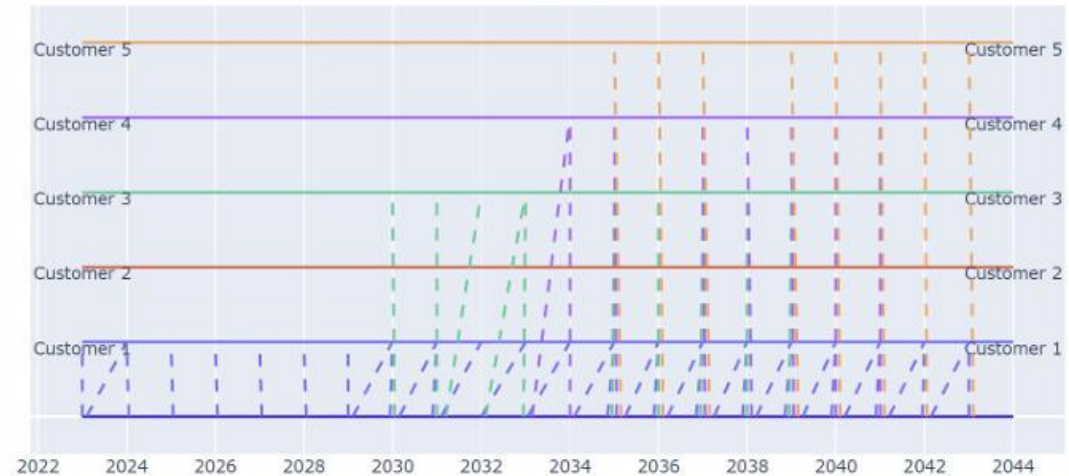
In MILP:

- 378 integer variables
- 108 equality constraints
- 111 inequality constraints

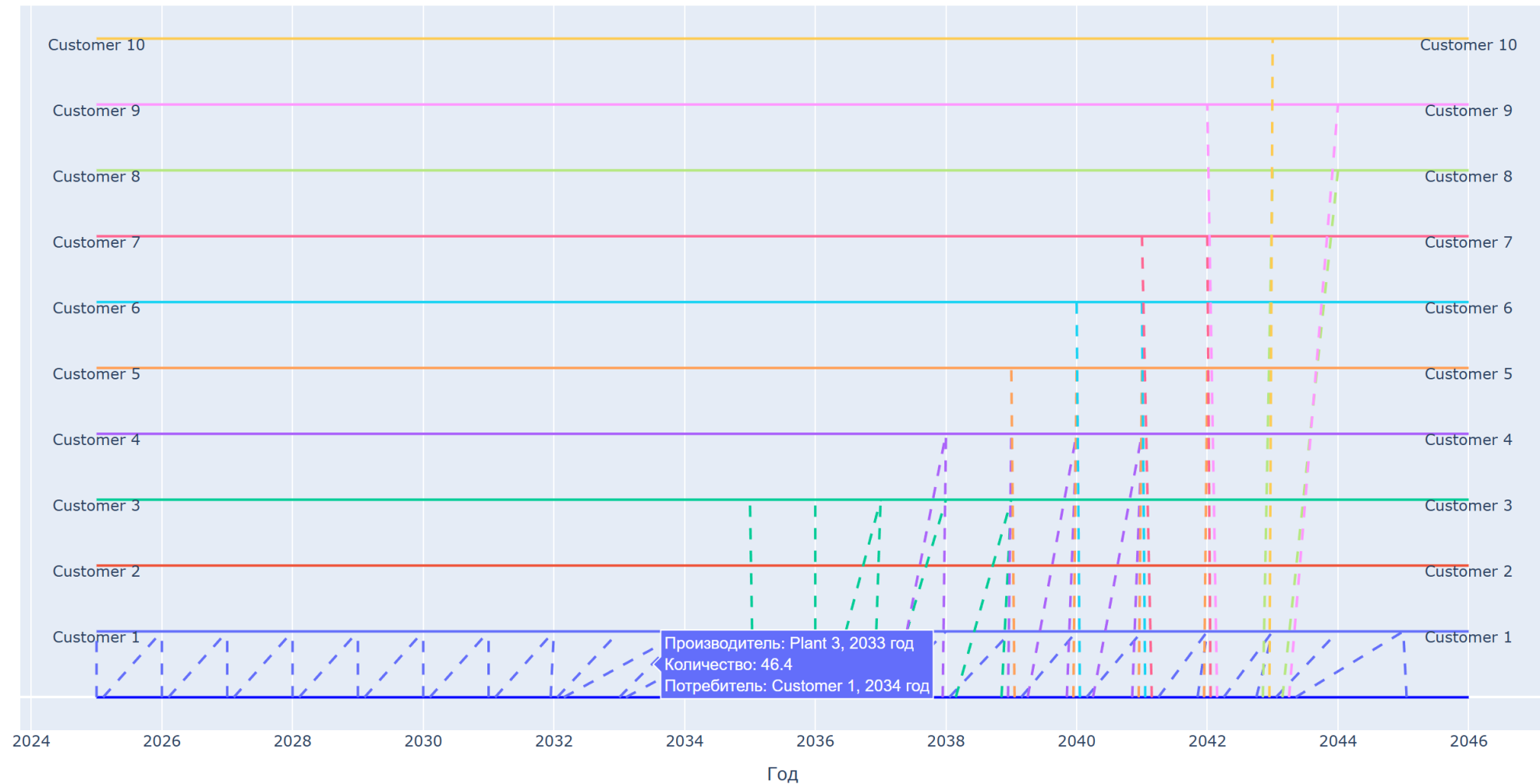
In QUBO:

- 1233 binary decision variables

8 hours of manual planning to
5 min quantum algorithm run



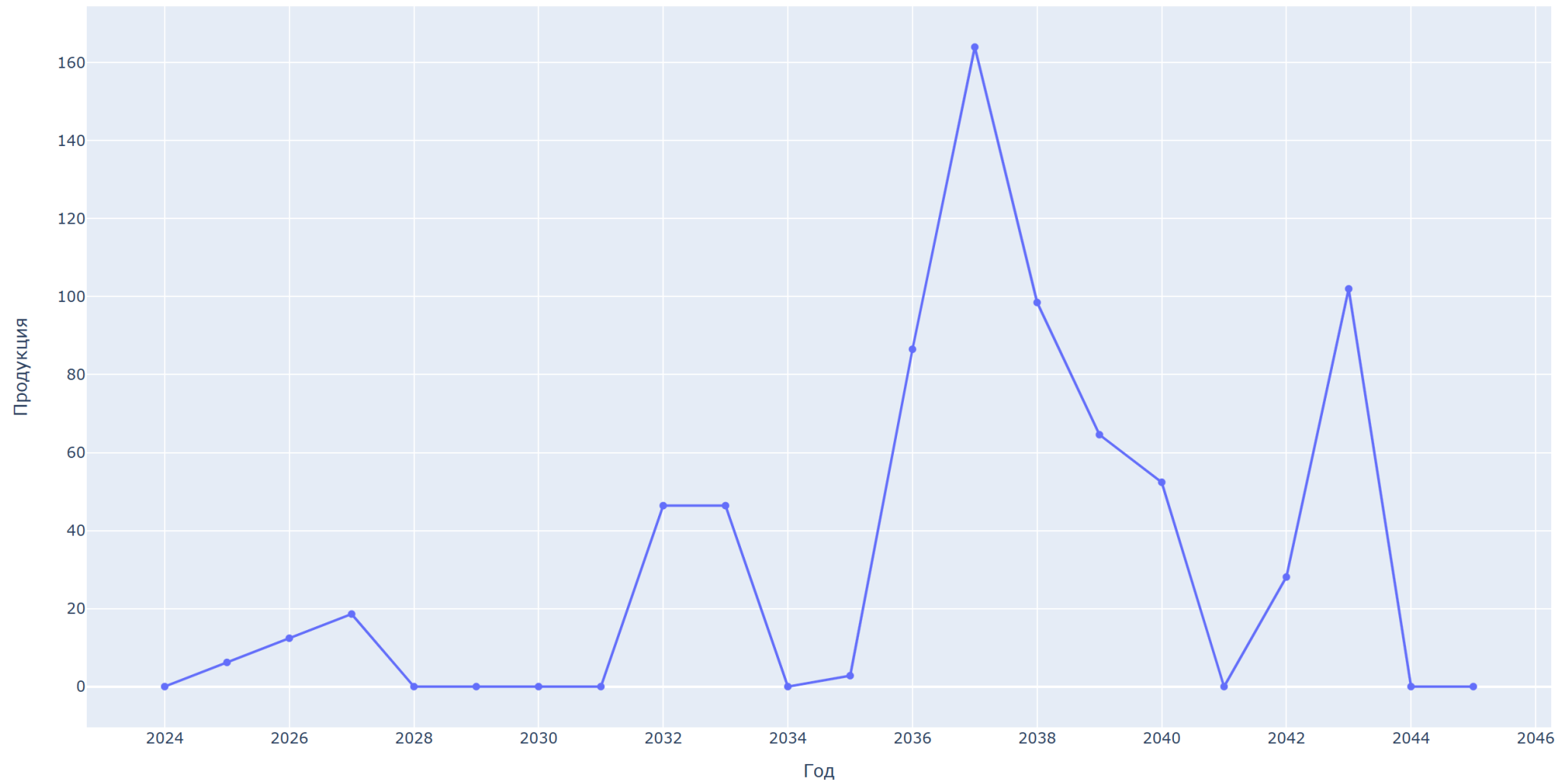
Production plan



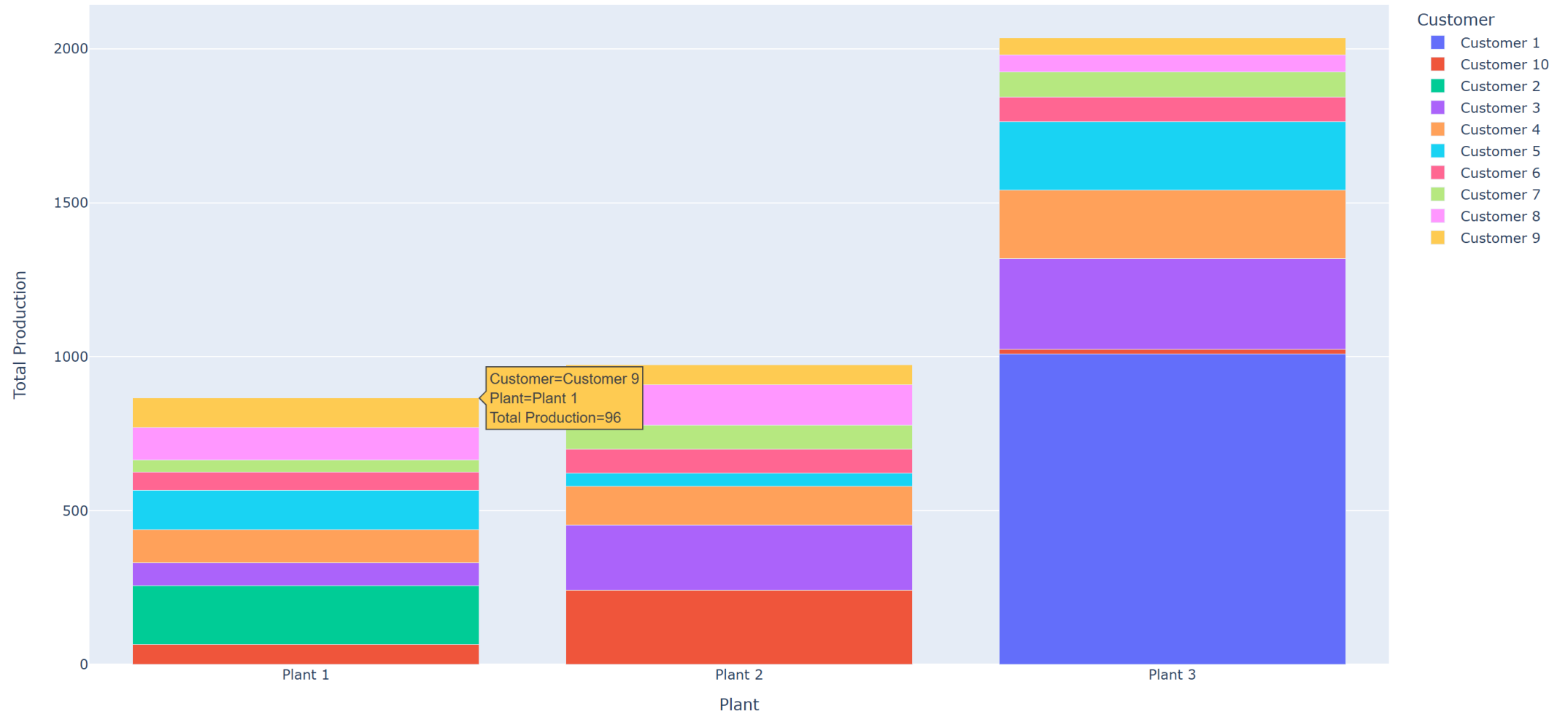
Customer consumption



Plant storage



Production distribution



Results

- Provided verification of quantum inspired algorithms for real world problems without decomposition
 - In the literature often small test problems are solved.
- Solved the problem on real data
 - Quantum algorithms provided several solutions which could be used for stability analysis
- Rather small problem size for testing on real quantum computer
 - In the future could be run on real machine
- Comparison with HIGHS solver was provided
 - Results are identical

Our plan:

- New model with additional constraints was developed
- New real data with bigger scale obtained

Thanks for your attention!

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