

# An Introduction to Quantum Computing and Quantum Machine Learning with Quantum Parametric Circuits



Ilya Sinayskiy

24 June 2025



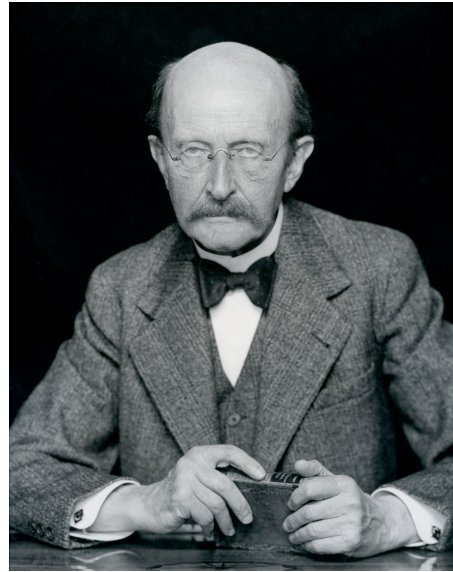
# Outline

- History of Quantum Computing
- Universal error correcting fault -tolerant quantum computers vs NISQ devices
- Basic of Quantum Computations
- QPC examples VQE/VQLS/QAOA/QCNN
- Application of architectural search to Ansatz design



# Quantum Computers

Quantum Mechanics:  
early 1900



Computer Science:  
1930



Richard P. Feynman



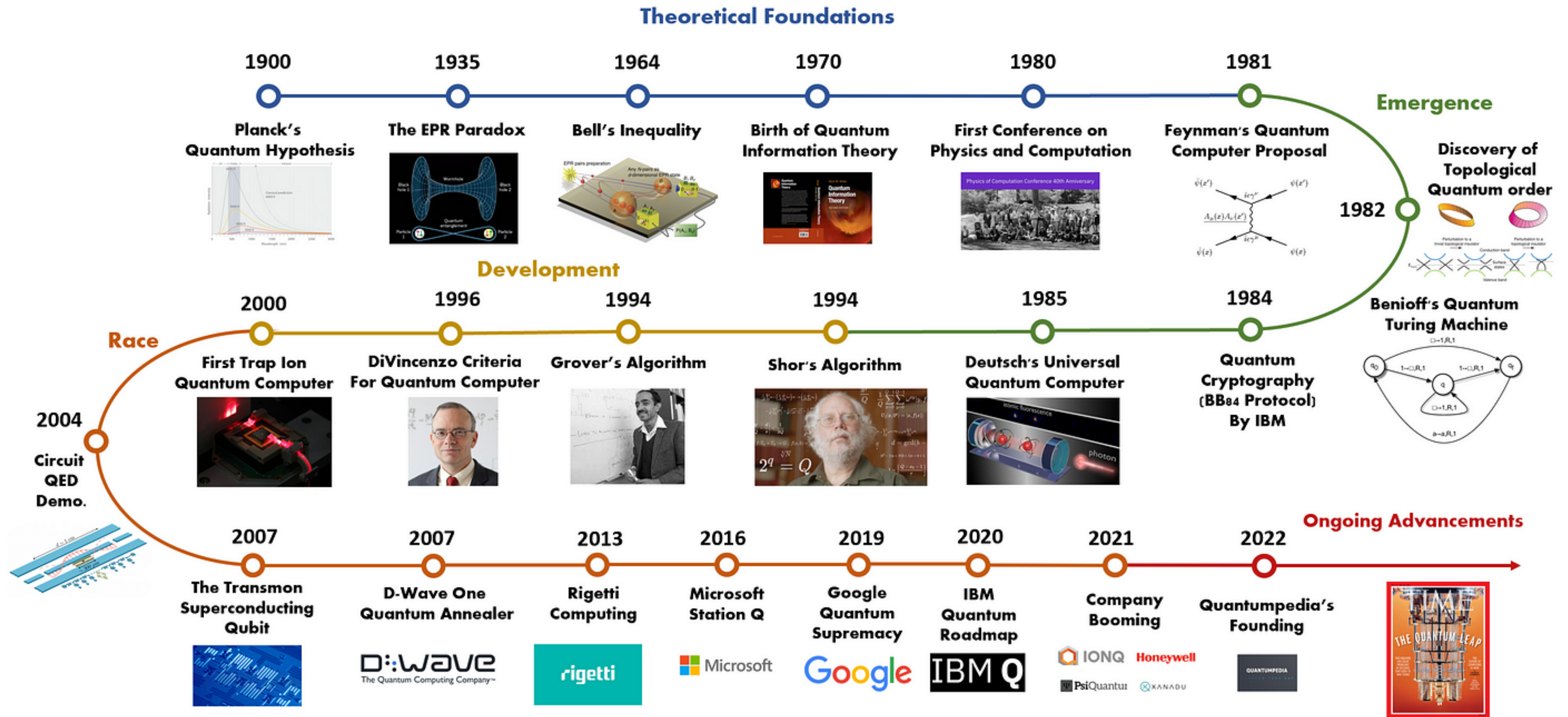
Yuri I. Manin (Юрий Манин)

Yu. I. Manin, “Computable and Non-Computable,” Sovetskoe Radio, Moscow, 1980, 128 p.

R.P. Feynman, Simulating physics with computers, Internat. J. Theoret. Phys. 21 (6–7) (1982) 467–488, <http://dx.doi.org/10.1007/bf02650179>.



# Quantum Computers



<https://quantumpedia.uk/a-brief-history-of-quantum-computing-e0bbd05893d0>

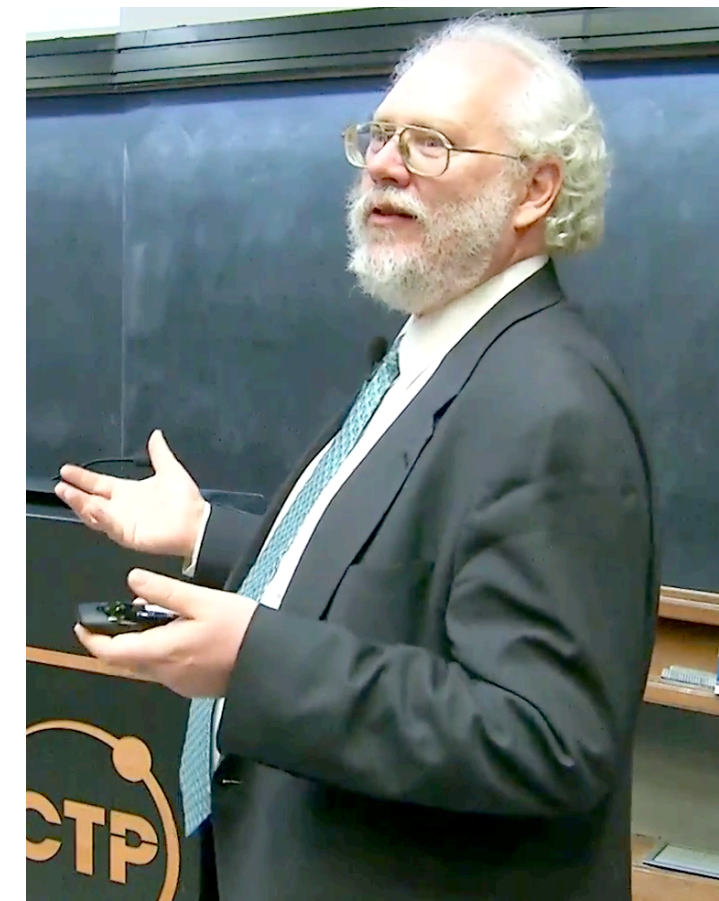


# Quantum Computers

## Shor's Algorithm

*Shor, P.W. (1994). "Algorithms for quantum computation: Discrete logarithms and factoring". Proceedings 35th Annual Symposium on Foundations of Computer Science. pp. 124–134. doi:10.1109/sfcs.1994.365700.*

*Shor, Peter W. (October 1997). "Polynomial-Time Algorithms for Prime Factorization and Discrete Logarithms on a Quantum Computer". SLAM Journal on Computing. 26 (5): 1484–1509. arXiv:quant-ph/9508027. doi:10.1137/S0097539795293172. 2337707.*



## Shor's Code - QEC

*Shor, Peter W. (1995). "Scheme for reducing decoherence in quantum computer memory". Physical Review A. 52 (4): R2493 – R2496*

## Fault-tolerant quantum computation

*Shor, Peter W. (1997). "Fault-tolerant quantum computation" quant-ph/9605011; 37th Symposium on Foundations of Computing, IEEE Computer Society Press, 1996, pp. 56–65*

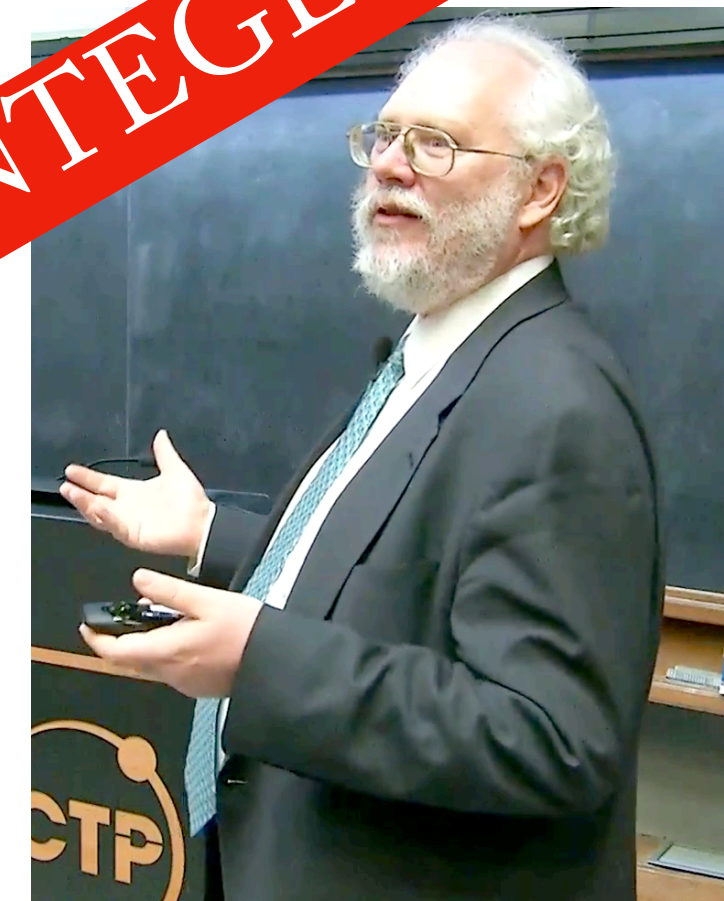


# Quantum Computers

## Shor's Algorithm

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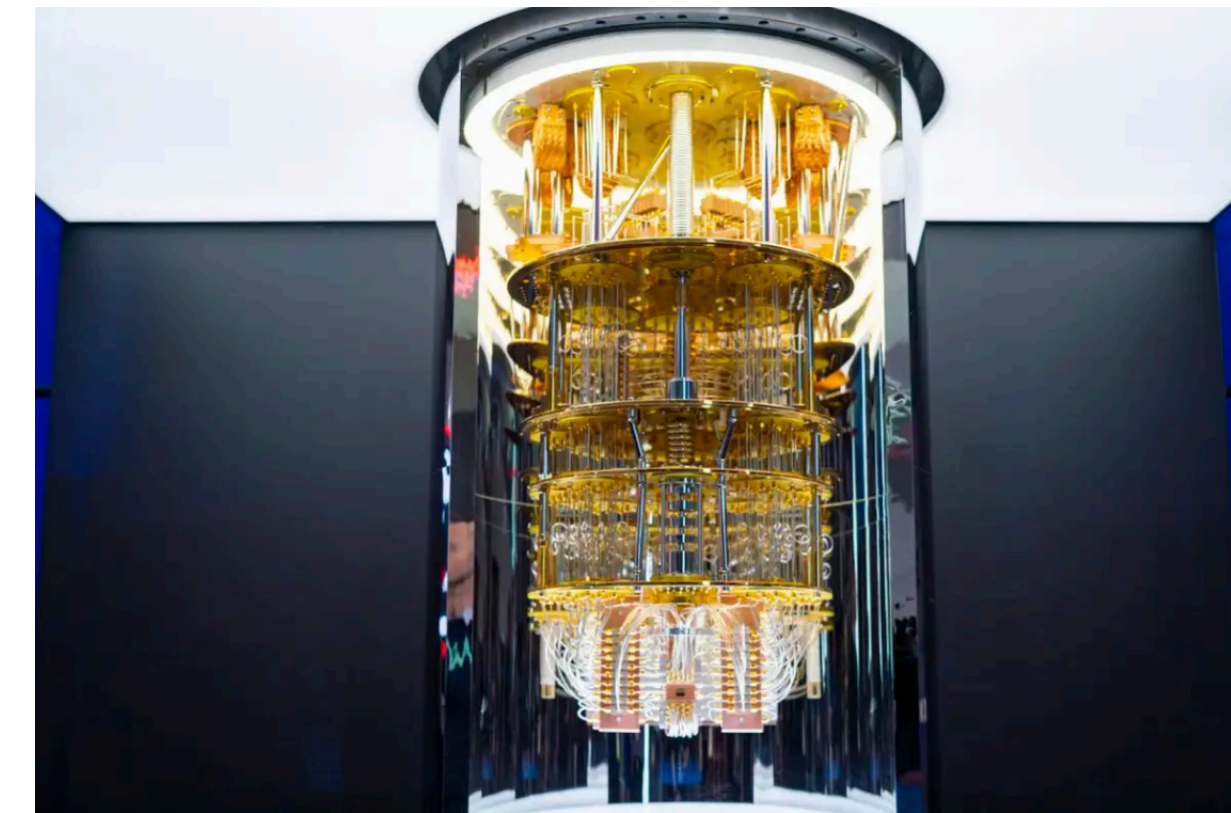
# Quantum Computers - Current Status (NISQ)

## Google - Willow Chip (2024)



105 qubit + some QEC

## IBM - Condor Chip (2025)



1121 qubit



# Quantum Computers - Current Status

## Google - Willow Chip (2024)



105 qubit + QEC

## IBM - Condor Chip (2025)



1121 qubit

FAR AWAY FROM FACTORING 2048bit INTEGER



# Math of Quantum Computing on a single page

- **Qubit** is superposition of 0 and 1:  $\alpha_0 |0\rangle + \alpha_1 |1\rangle \in \mathbb{C}^2$
- $n$ -qubit system: superposition of all  $n$ -bit strings:

$$\sum_{x \in \{0,1\}^n} \alpha_x |x\rangle \in \mathbb{C}^{2^n}$$

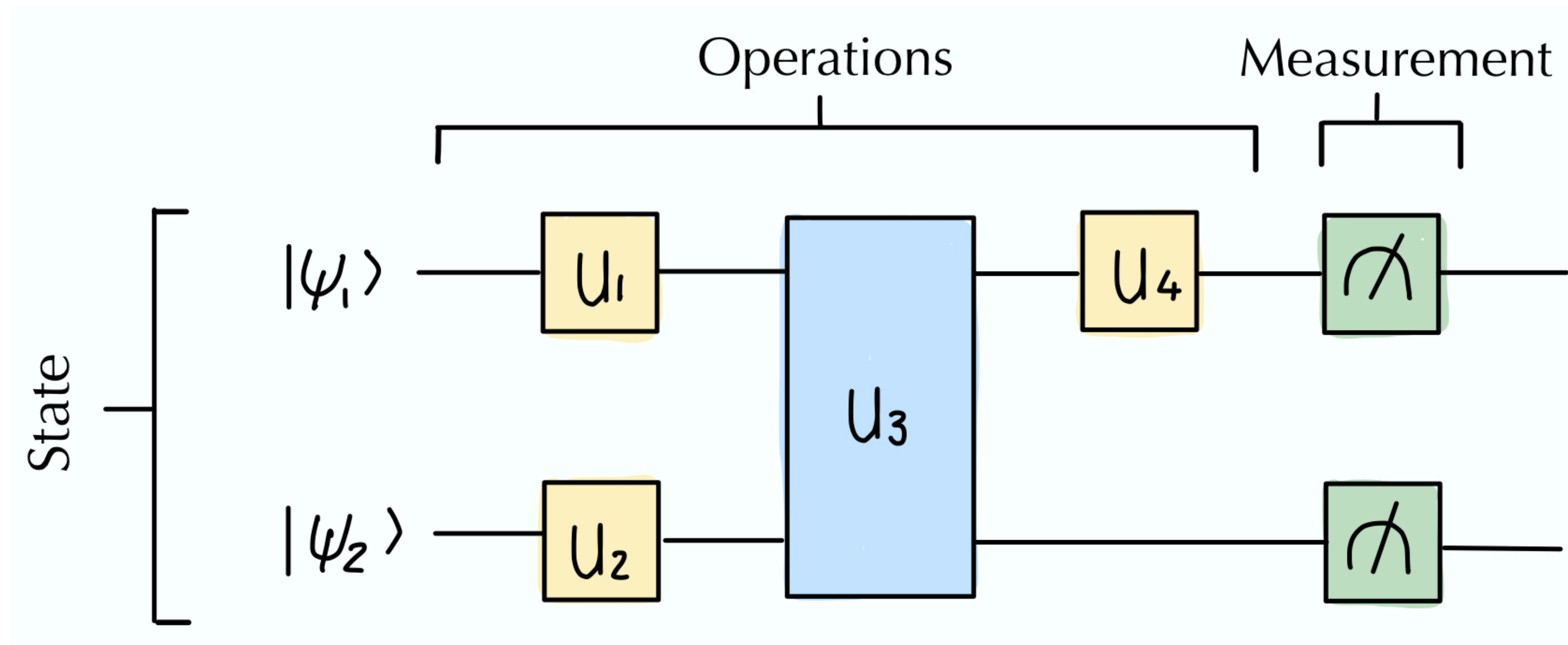
- **Measurement**: see outcome  $x \in \{0,1\}^n$  with probability  $|\alpha_x|^2$
- **Unitary transformation**: matrix that preserves the length of the vector of amplitudes. **Gates**: unitaries on 1 qubit

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}, \quad H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

- or on 2 qubits, CNOT:  $|a, b\rangle \mapsto |a, a \oplus b\rangle$   $CNOT = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

- Combine simultaneous gates via tensor product, combine sequential gates via matrix product

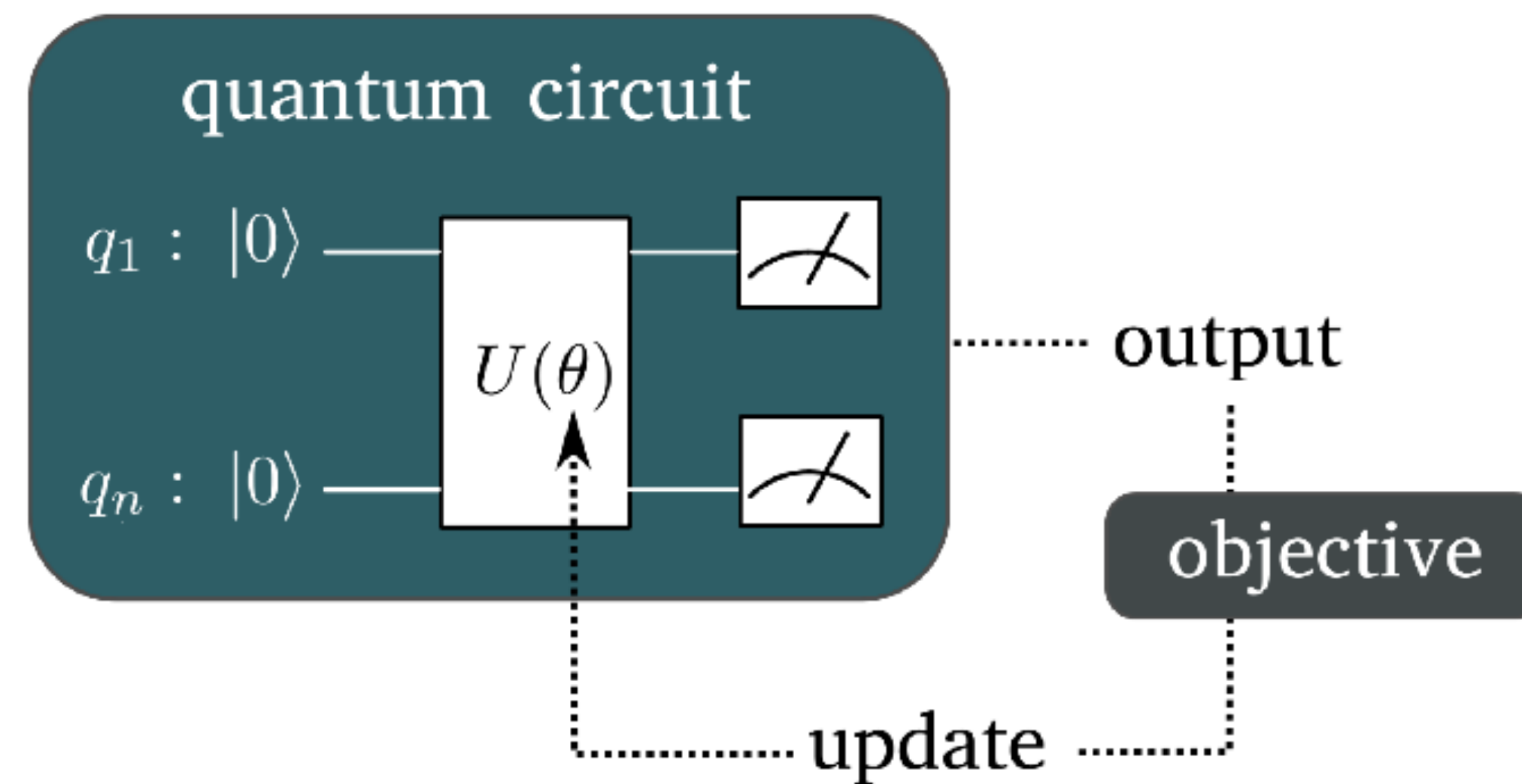
# Circuit Model of Quantum Computations



Circuit for Quantum Computer

# Quantum Computing with NISQ devices

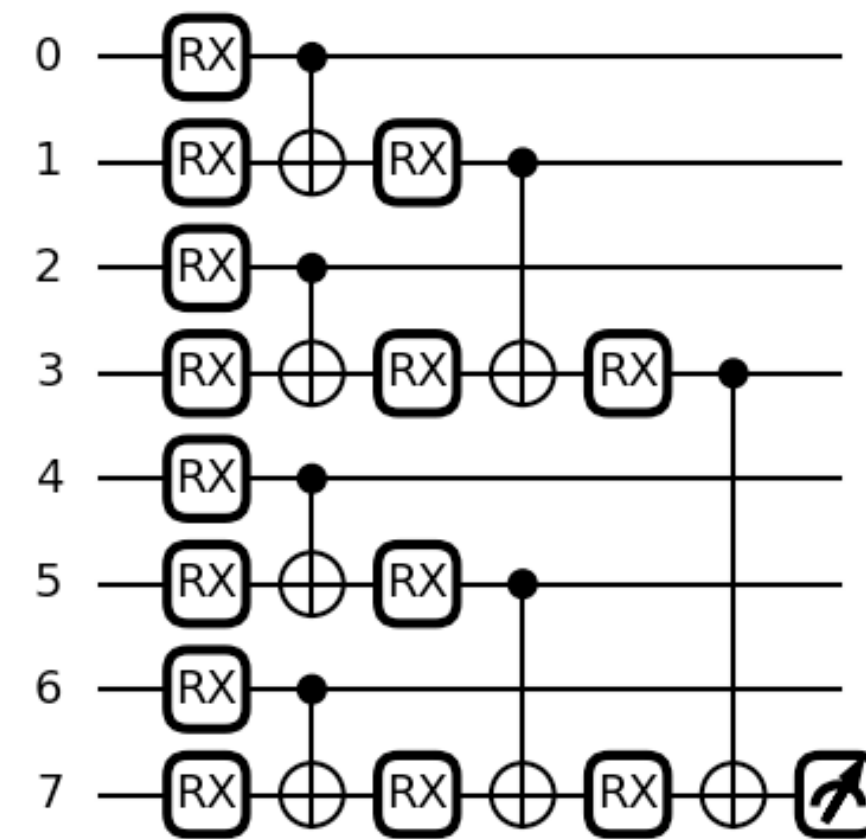
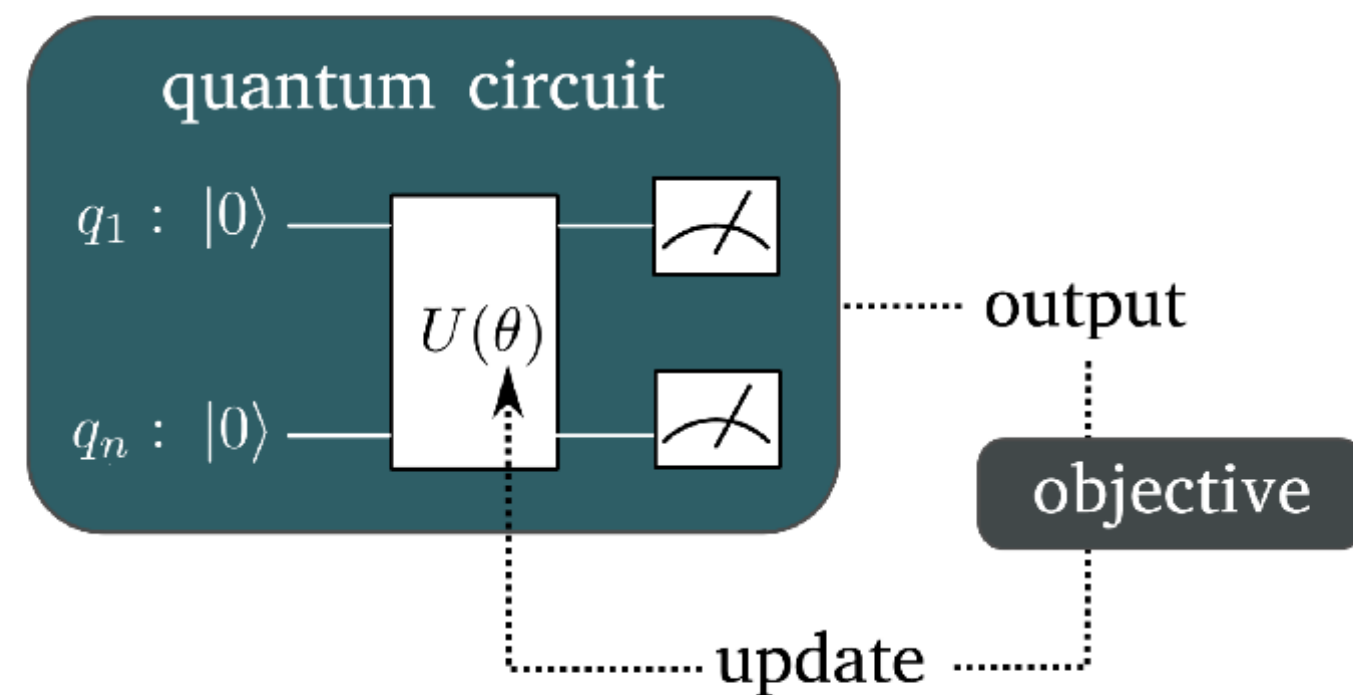
## Quantum Parametric Circuit





# Quantum Computing with NISQ devices

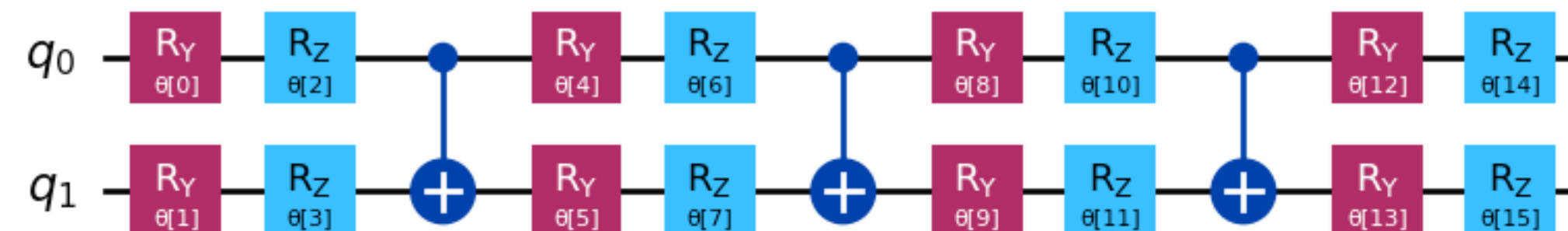
## Quantum Parametric Circuit



$$RX(\theta) = \exp\left(-i\frac{\theta}{2}X\right) = \begin{pmatrix} \cos\left(\frac{\theta}{2}\right) & -i\sin\left(\frac{\theta}{2}\right) \\ -i\sin\left(\frac{\theta}{2}\right) & \cos\left(\frac{\theta}{2}\right) \end{pmatrix}$$

$$RZ(\phi) = \exp\left(-i\frac{\phi}{2}Z\right) = \begin{pmatrix} e^{-i\frac{\phi}{2}} & 0 \\ 0 & e^{i\frac{\phi}{2}} \end{pmatrix}$$

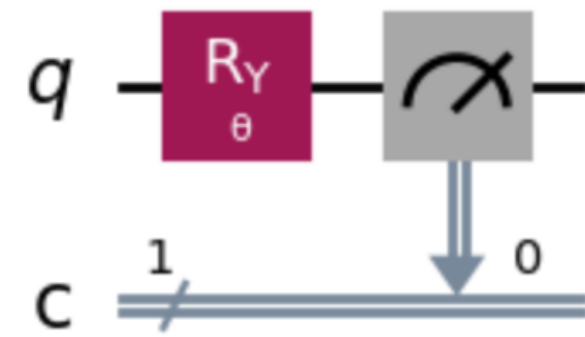
$$RY(\theta) = \exp\left(-i\frac{\theta}{2}Y\right) = \begin{pmatrix} \cos\left(\frac{\theta}{2}\right) & -\sin\left(\frac{\theta}{2}\right) \\ \sin\left(\frac{\theta}{2}\right) & \cos\left(\frac{\theta}{2}\right) \end{pmatrix}$$



# QPC for Ground State of a Two Level System

$$\hat{H}_{\text{TLS}} = \omega \hat{Z} + \lambda \hat{X}$$

$$\text{Ansatz } |\psi(\theta)\rangle = \hat{R}_y(\theta)|0\rangle$$



$$\langle\psi(\theta)|\hat{Z}|\psi(\theta)\rangle = \cos\theta$$



$$\langle\psi(\theta)|\hat{H}\hat{Z}\hat{H}|\psi(\theta)\rangle = \langle\psi(\theta)|\hat{X}|\psi(\theta)\rangle = \sin\theta$$

$$\langle\psi(\theta)|\hat{H}_{\text{TLS}}|\psi(\theta)\rangle = \omega \cos\theta + \lambda \sin\theta$$

$$\min_{\theta}(\langle\psi(\theta)|\hat{H}_{\text{TLS}}|\psi(\theta)\rangle) = -\sqrt{\omega^2 + \lambda^2}$$

# General case: QPC for Ground State Estimation

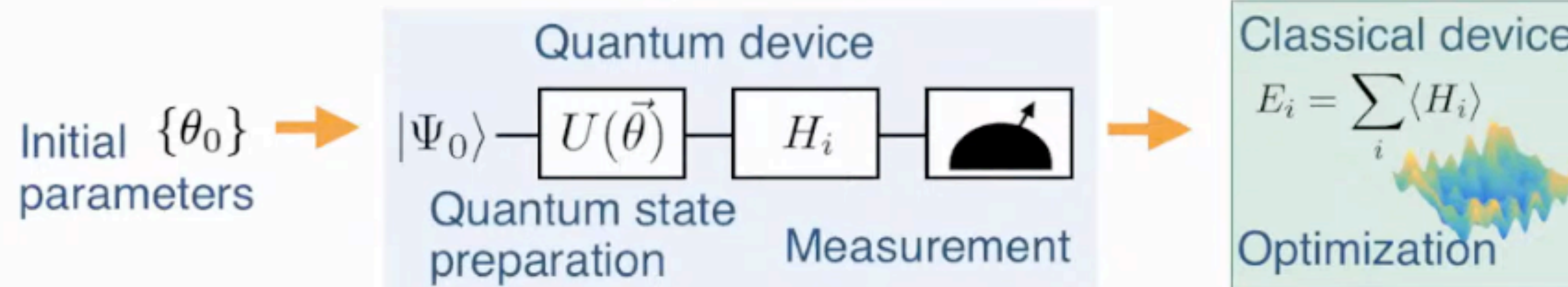
## The Variational Quantum Eigensolver

1. **Quantum** state preparation (use parameteric ansatz to describe the wavefunction)
2. **Classical** Optimization (variational principle)

$$\Psi(\vec{\theta}) = U(\vec{\theta})|\Psi_0\rangle$$

$$E = \min_{\vec{\theta}} \langle \Psi(\vec{\theta}) | \hat{H} | \Psi(\vec{\theta}) \rangle \geq E_{\text{exact}}$$

$\{\theta_1\}$





# Solving Systems of Linear Equations with QPC

PRL **103**, 150502 (2009)

PHYSICAL REVIEW LETTERS

week ending  
9 OCTOBER 2009



## Quantum Algorithm for Linear Systems of Equations

Aram W. Harrow,<sup>1</sup> Avinatan Hassidim,<sup>2</sup> and Seth Lloyd<sup>3</sup>

<sup>1</sup>*Department of Mathematics, University of Bristol, Bristol, BS8 1TW, United Kingdom*

<sup>2</sup>*Research Laboratory for Electronics, MIT, Cambridge, Massachusetts 02139, USA*

<sup>3</sup>*Research Laboratory for Electronics and Department of Mechanical Engineering, MIT, Cambridge, Massachusetts 02139, USA*

(Received 5 July 2009; published 7 October 2009)

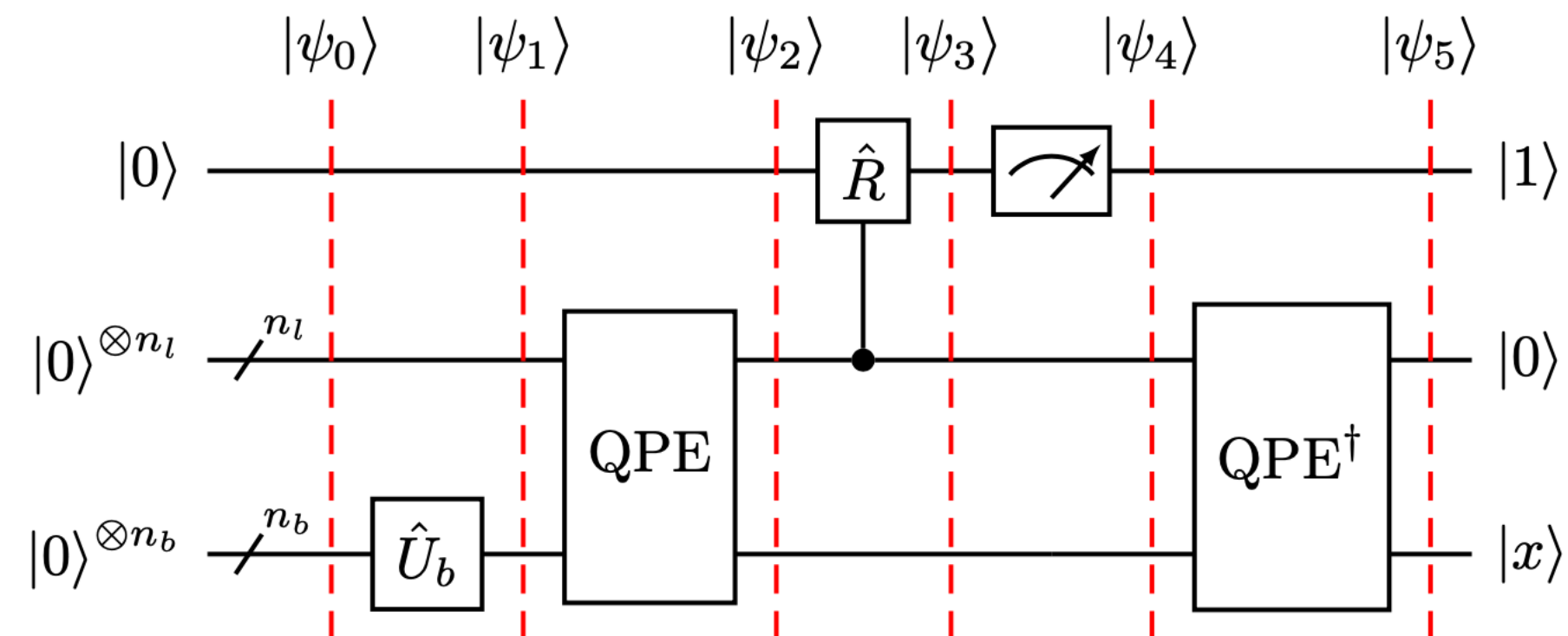
Solving linear systems of equations is a common problem that arises both on its own and as a subroutine in more complex problems: given a matrix  $A$  and a vector  $\vec{b}$ , find a vector  $\vec{x}$  such that  $A\vec{x} = \vec{b}$ . We consider

Classical algorithms scales as  $N\sqrt{\kappa}$ .

HHL scales as  $\text{poly}(\log(N), \kappa)$

$$A\vec{x} = \vec{b}. \longrightarrow A|x\rangle = |b\rangle, \xrightarrow{U = e^{iAt}} |x\rangle = A^{-1}|b\rangle.$$

HHL works in fault-tolerant settings only!



# Solving Systems of Linear Equations with QPC

$$|x\rangle := \frac{\sum_i x_i |i\rangle}{\|\sum_i x_i |i\rangle\|_2},$$

$$|b\rangle := \frac{\sum_i b_i |i\rangle}{\|\sum_i b_i |i\rangle\|_2}.$$

$$C_G = \frac{\langle x | H_G | x \rangle}{\langle \psi | \psi \rangle},$$

$$H_G = A^\dagger (\mathbb{1} - |b\rangle\langle b|) A.$$

Global CF is susceptible to BP!

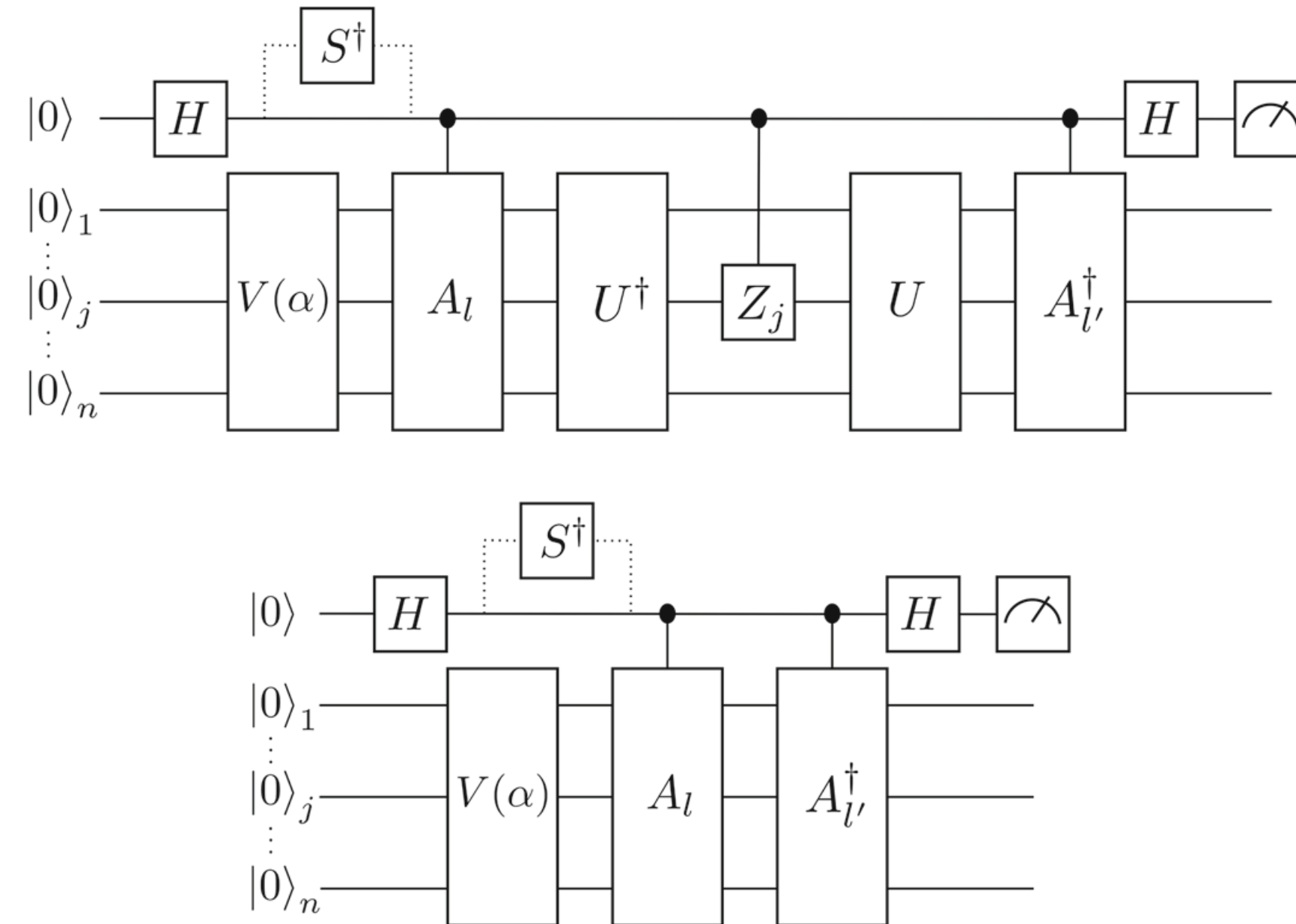
$$H_L = A^\dagger U \left( \mathbb{1} - \frac{1}{n} \sum_{j=1}^n |0_j\rangle\langle 0_j| \otimes \mathbb{1}_{\bar{j}} \right) U^\dagger A,$$

$$C_L = \frac{\langle x | H_L | x \rangle}{\langle \psi | \psi \rangle}.$$

$$\mathbf{A} = \sum_{i=1}^m c_i A_i$$

$$\tilde{\mathbf{A}} = \begin{pmatrix} 0 & \mathbf{A} \\ \mathbf{A}^\dagger & 0 \end{pmatrix}$$

$$|\psi\rangle = A|x\rangle \quad U|0\rangle = |b\rangle$$



**Fig. 2** Hadamard Test Circuits for cost function  $C_L$  Eq. (9). The top circuit is employed when calculating the value of the numerator  $\langle x | H_L | x \rangle$ , while the bottom circuit is employed when calculating the value of the denominator,  $\langle \psi | \psi \rangle$ . The  $S^\dagger$  gate is included when calculating imaginary-valued parts of the cost function and excluded when calculating the real-valued parts, and therefore, drawn in as a dotted line.  $V(\alpha)$  denotes the ansatz,  $A_l$  the  $l$ -th unitary from the linear sum  $A = \sum_i c_i A_i$  and  $U$  the unitary such that  $U|0\rangle = |b\rangle$ .  $Z_j$  denotes a standard  $Z$  gate on the  $j$ -th qubit, and  $H$  being a standard Hadamard gate

Carlos Bravo-Prieto, Ryan LaRose, M. Cerezo, Yigit Subasi, Lukasz Cincio, and Patrick J. Coles, **Variational Quantum Linear Solver**, Quantum 7, 1188 (2023).

Pellow-Jarman, A., Sinayskiy, I., Pillay, A. *et al.* Near term algorithms for linear systems of equations. *Quantum Inf Process* **22**, 258 (2023). <https://doi.org/10.1007/s11128-023-04020-2>

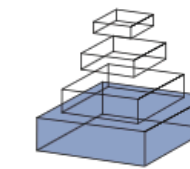
# Quantum Approximate Optimization Algorithm (QAOA)

Many problems in Optimization theory/Graph Theory/Computer Science can be formulated as a minimum energy state of classical spin glass.

frontiers in  
**PHYSICS**

**REVIEW ARTICLE**

published: 12 February 2014  
doi: 10.3389/fphy.2014.00005



## Ising formulations of many NP problems

**Andrew Lucas \***

*Lyman Laboratory of Physics, Department of Physics, Harvard University, Cambridge, MA, USA*

Partitioning problems (number partitioning, graph partitioning)  
Cliques, covering and packing problems (exact cover, set packing, vertex cover)  
Satisfiability, minimal maximal matching, knapsack with integer weights, coloring problems  
Hamiltonian cycles and paths, traveling salesman, tree problems...

# Quantum Approximate Optimization Algorithm (QAOA)

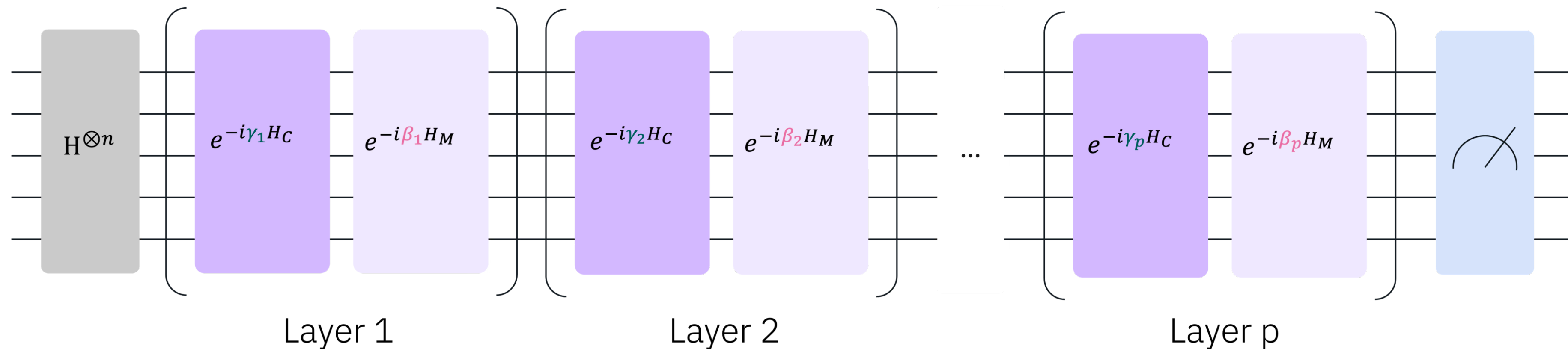
$$H(t) = \left(1 - \frac{t}{T}\right) H_0 + \frac{t}{T} H_P,$$

$$H_0 = -h_0 \sum_{i=1}^N \sigma_i^x,$$

$$H(s_1, \dots, s_N) = - \sum_{i < j} J_{ij} s_i s_j - \sum_{i=1}^N h_i s_i.$$

We go from the ground state  
of the driver's Hamiltonian

To the ground state of  
the problem Hamiltonian





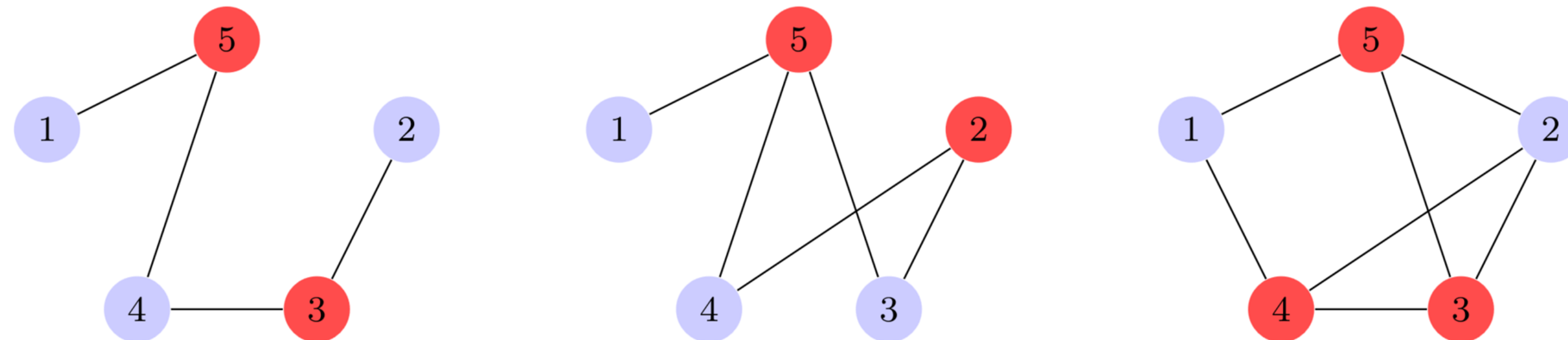
# QAOA example: Minimum Vertex Cover

$$\text{Minimize: } \sum_{i \in V} x_i$$

$$\text{Subject to: } x_i + x_j \geq 1, \quad \forall (i, j) \in E$$

$$\text{and: } x_i \in \{0, 1\}, \quad \forall i \in V$$

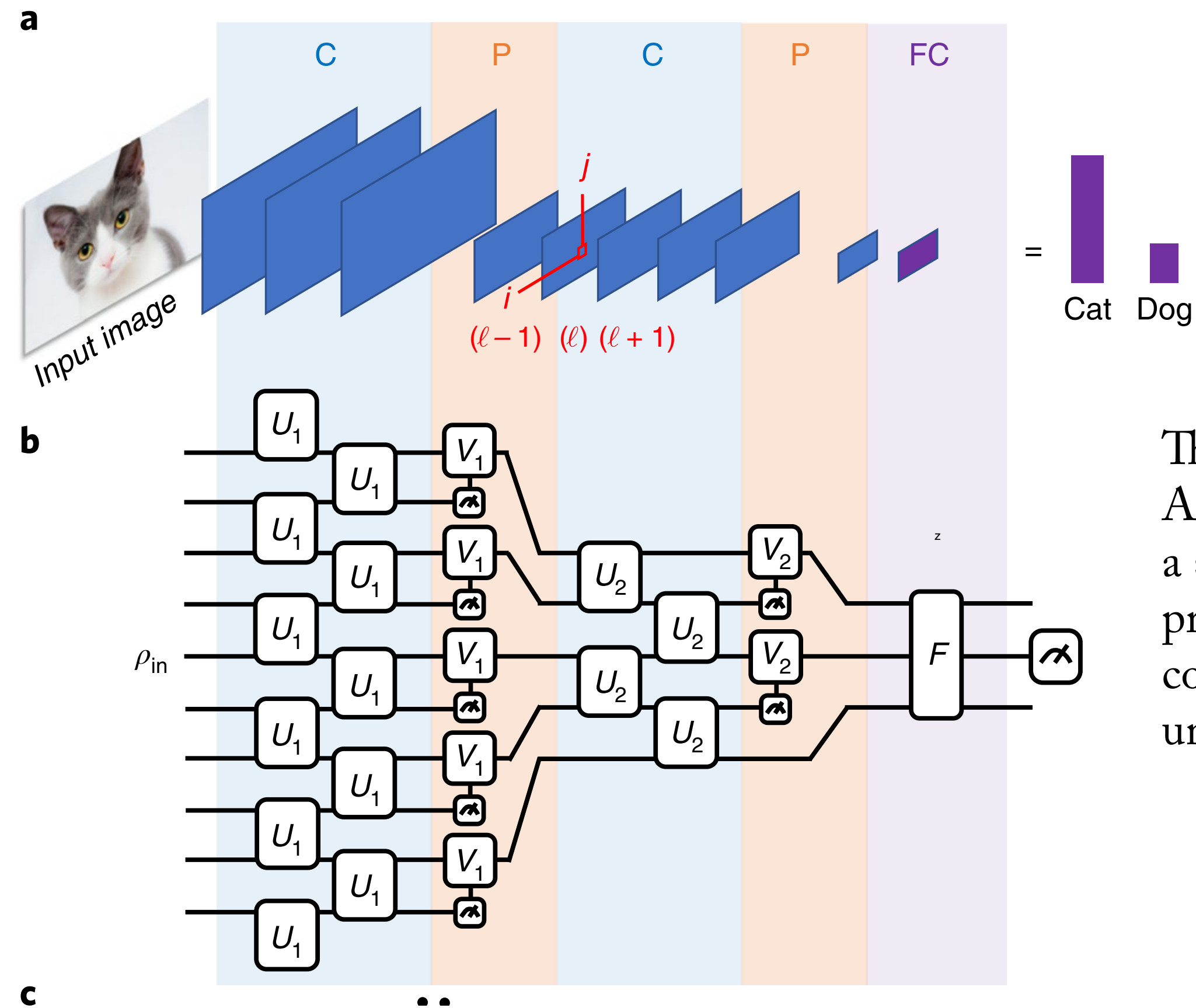
$$H_C = A \sum_{(u,v) \in E} (1 - x_u)(1 - x_v) + B \sum_{v \in V} x_v$$



**Figure 1.** In the three graphs above, the red nodes show the set of vertices forming each graph's respective minimum vertex cover. Each edge in the graph under consideration must have, at least one vertex in the cover. The cover forms the minimum cover of a graph, when it contains the fewest number of vertices, whilst ensuring each edge is still incident to at least one.

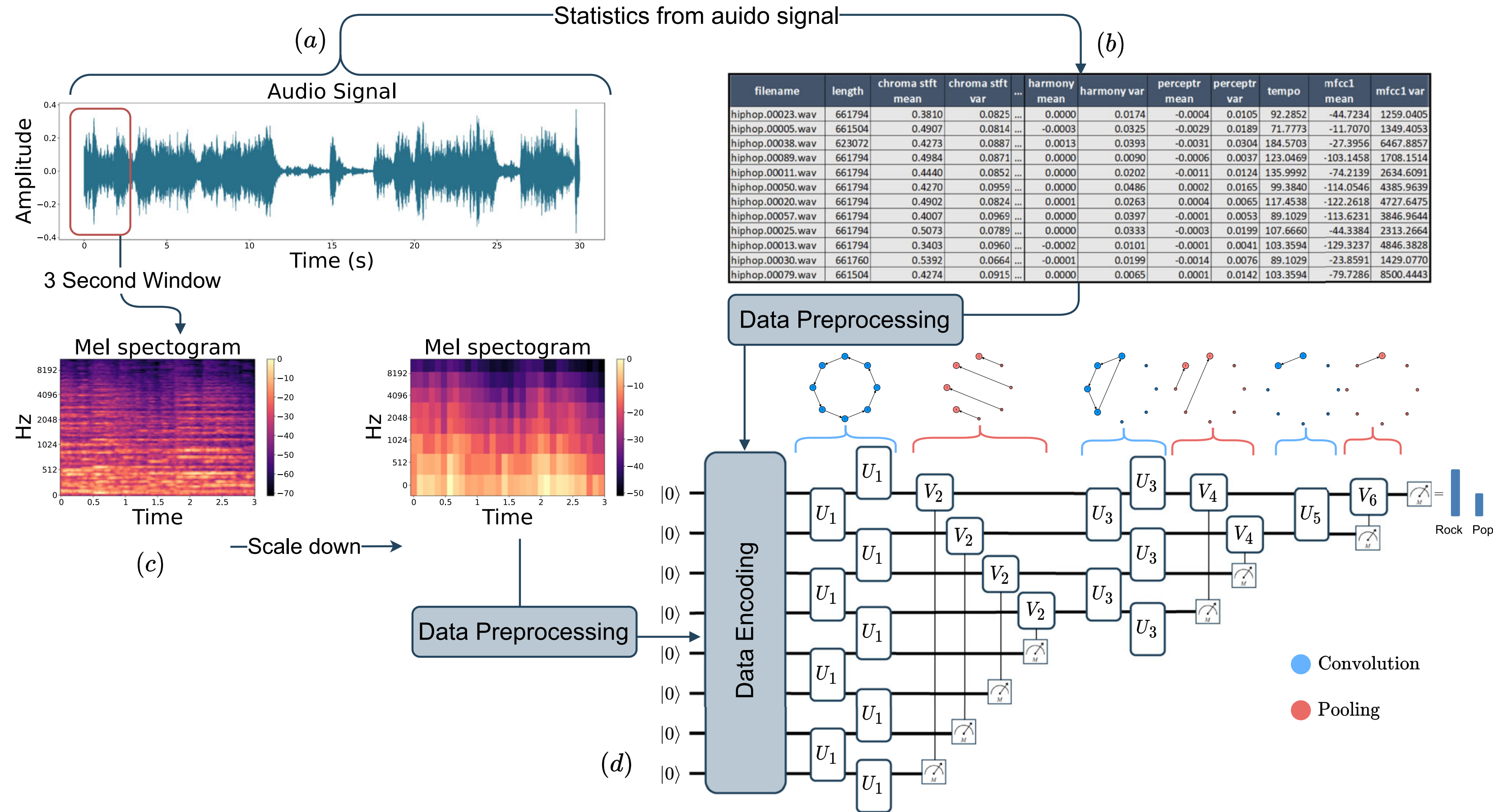


# Quantum Convolutional Neural Networks



The concept of QCNNs. (a) Simplified illustration of classical CNNs. A sequence of image-processing layers transforms an input image into a series of feature maps (blue rectangles) and finally into an output probability distribution (purple bars). C, convolution; P, pooling; FC, fully connected. (b) QCNNs inherit a similar layered structure. Boxes represent unitary gates or measurement with feed-forwarding.

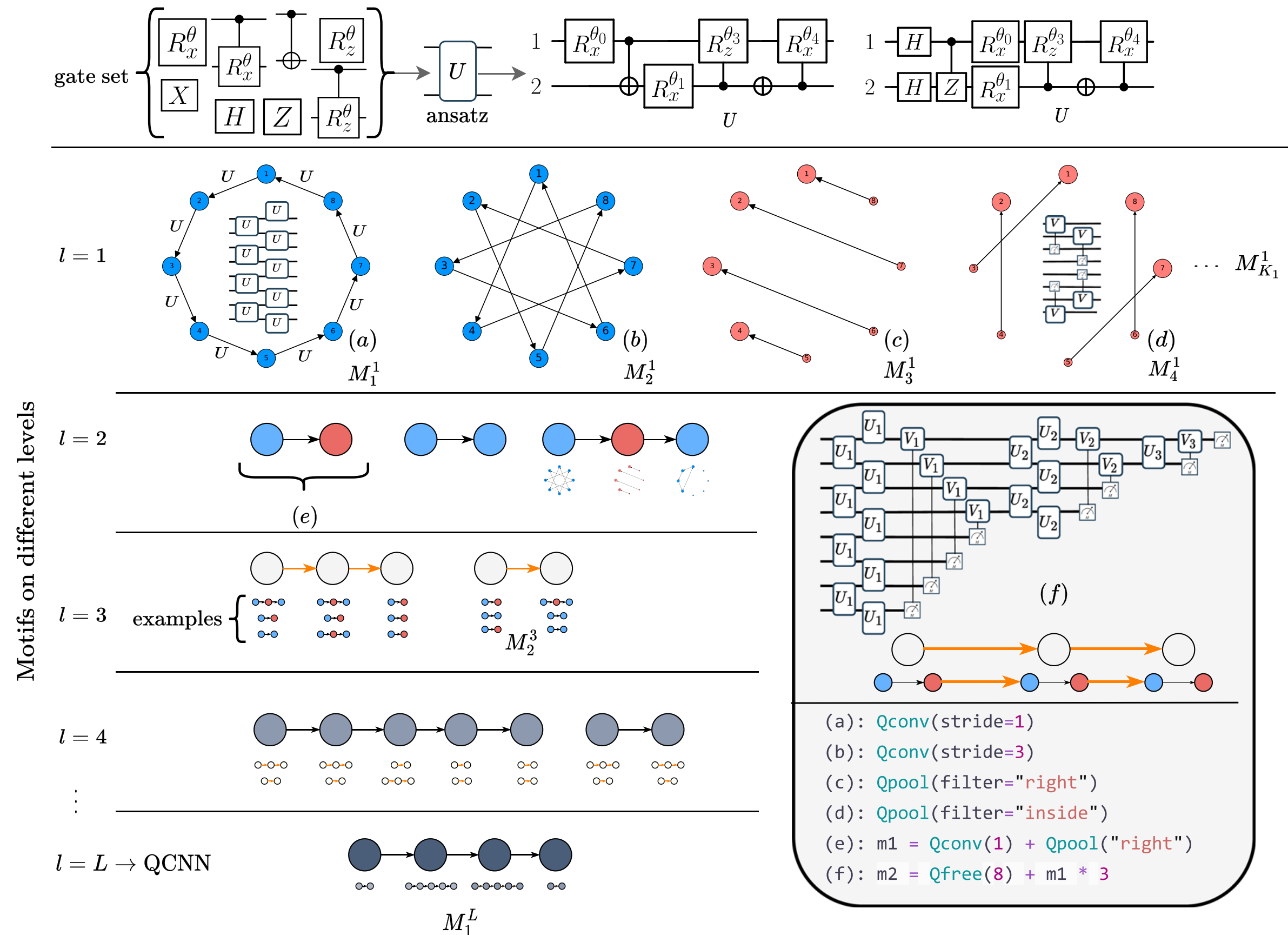
# Quantum Convolutional Neural Networks



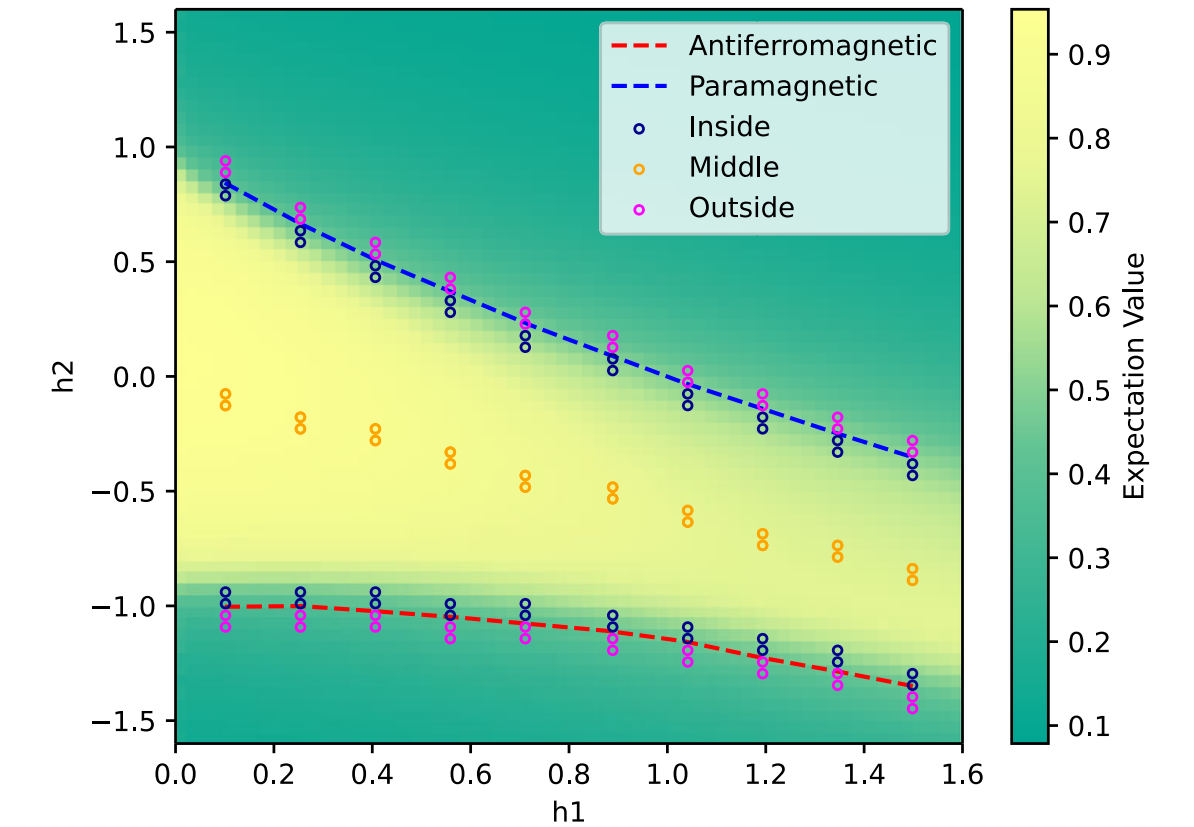
**Fig. 1 Machine-learning pipeline for music genre classification.** The machine-learning pipeline we implemented for music genre classification. Given an audio signal of a song (a), we generate two forms of data: tabular (b) and image (c). Each form has data preprocessing applied before being encoded into a quantum state (d). The QCNN circuit shown in (d) favours Principal Component Analysis (PCA) because qubits are pooled from bottom to top, and principal components are encoded from top to bottom. This architecture is an instance of the reverse binary tree family that we generated with our framework.



# Quantum Convolutional Neural Networks



**Fig. 2 Hierarchical quantum circuit representation.** An overview of our architectural representation for QCNNs. From a given set of gates, we build two-qubit unitary ansatzes. The representation then captures design motifs  $M_k^l$  on different levels  $l$  of the hierarchy. On the lowest level  $l=1$ , we define primitives which act as building blocks for the architecture. For example, a convolution operation with stride one is encoded as the directed graph  $M_1^1$ . The directed graph  $M_3^1$  is a pooling operation that measures the bottom half of the circuit. Combined, they form the level two motif (e): a convolution-pooling unit  $M_2^2$ . Higher-level motifs consist of combinations of lower-level motifs up until the final level  $l=L$ , which contains only one motif  $M_1^L$ , the complete QCNN architecture.  $M_1^L$  is a hierarchy of directed graphs fully specifying how to spread the unitary ansatzes across the circuit. The two lines of code (e) and (f) show the power of this representation as it is all that is required to create the entire QCNN circuit from Fig. 1d. The code comes from the Python package we implemented based on the work of this paper. It facilitates dynamic QCNN creation and search space design.



**Fig. 5 Quantum phase recognition result.** Expectation values for the circuit found via evolutionary search for a system of  $N=15$  spins. Points represent a test set of  $64 \times 64$  ground states for various  $h_1$  and  $h_2$  values of the Hamiltonian,  $J=1$ . The inside, middle and outside points were used to evaluate an architecture's fitness during search. The same colour scale as in<sup>5</sup> is used to facilitate comparison.

<b>Table 3.</b> Performance of architecture found with an evolutionary search.		
Metric	Reference	Found
Number of parameters	1308	<b>11</b>
Sample complexity (inside)	61.523	<b>36.079</b>
Sample complexity (middle)	<b>10.992</b>	13.253
MSE (outside)	<b>0.164</b>	0.167

Different performance metrics (lower is better) for the 15-qubit QCNN from ref. <sup>5</sup> and the architecture found via evolutionary search. The best performing architecture for each metric is highlighted in bold. Sample complexity represents the expected number of measurements required to be 95% confident that the ground state is in the SPT phase (non-zero expectation value). Metrics are calculated on a set of points in the test set, where inside refers to SPT points near the phase boundary, outside to non-SPT points near the phase boundary and middle to points in between, as shown in Fig. 5.

# Surprising Application: Ansatz Search

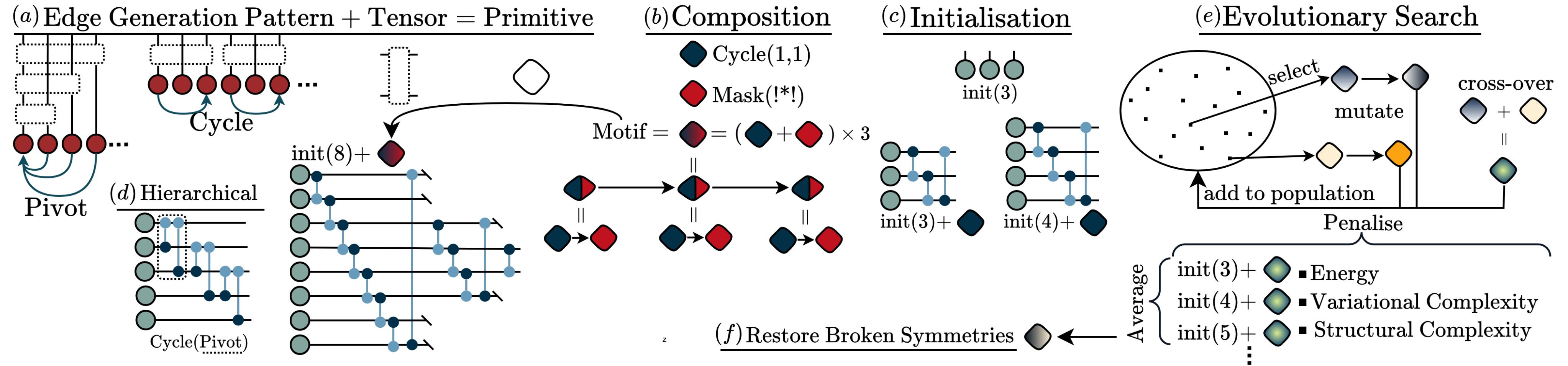


FIG. 1. Overview of our method, a domain-specific-language (DSL) enables ansatz generation via an evolutionary algorithm. (a) A primitive is an edge generation pattern associated with a tensor. (b) Composition: Sequences of primitives form motifs; sequences of motifs form higher-level motifs. (c) Specifying the number of nodes generates edges, and the associated tensor is repeated and connected to each edge, forming a tensor network. (d) A specified network, being itself a tensor, can again be associated with an edge generation pattern to form a new primitive. (e) The evolutionary algorithm mutates and crosses over motifs each generation. (f) Once the ansatz is found, broken symmetries are restored.



# Surprising Application: Ansatz Search

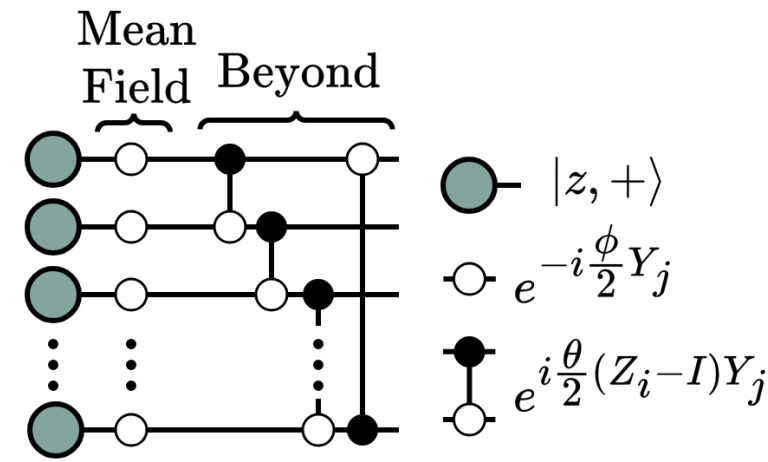


FIG. 2. The ansatz generated by our method for the LMG and TFIM models.

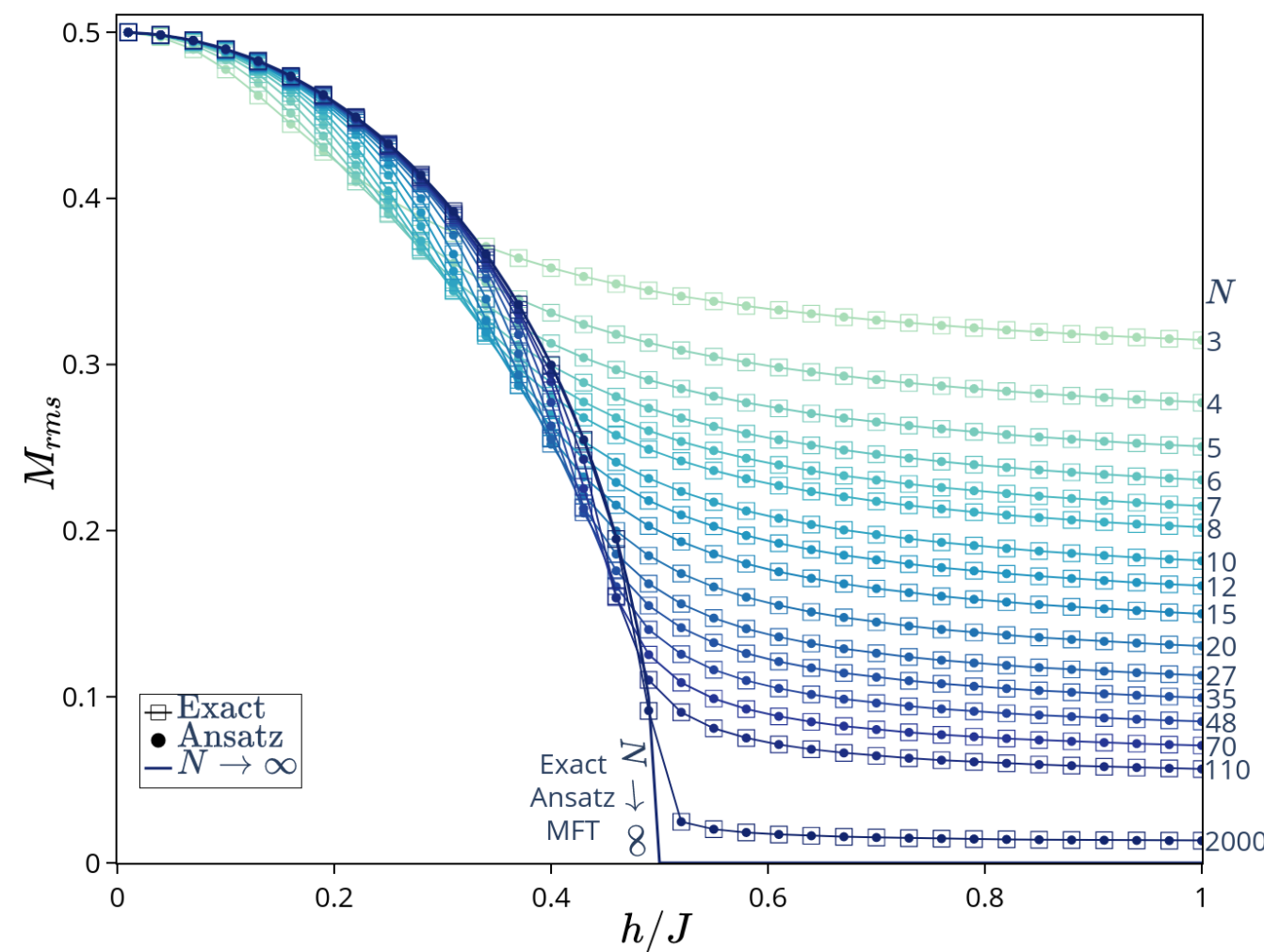


FIG. 3. RMS magnetisation of Eq. (20) vs  $h/J$  for the LMG model. Exact results compared to the symmetrised ansatz.

$$|\theta, \phi\rangle = \left( \prod_{k=0}^{N-1} C_{k,k+1}^{\theta} R_{k+1}^{\theta} \right) \left( \prod_{j=0}^{N-1} R_j^{\phi} \right) |z, +\rangle^{\otimes N}$$

$$C_{ij}^{\theta} = e^{i\frac{\theta}{2}Z_i Y_j} \quad \text{and} \quad R_j^{\theta} = e^{-i\frac{\theta}{2}Y_j}.$$

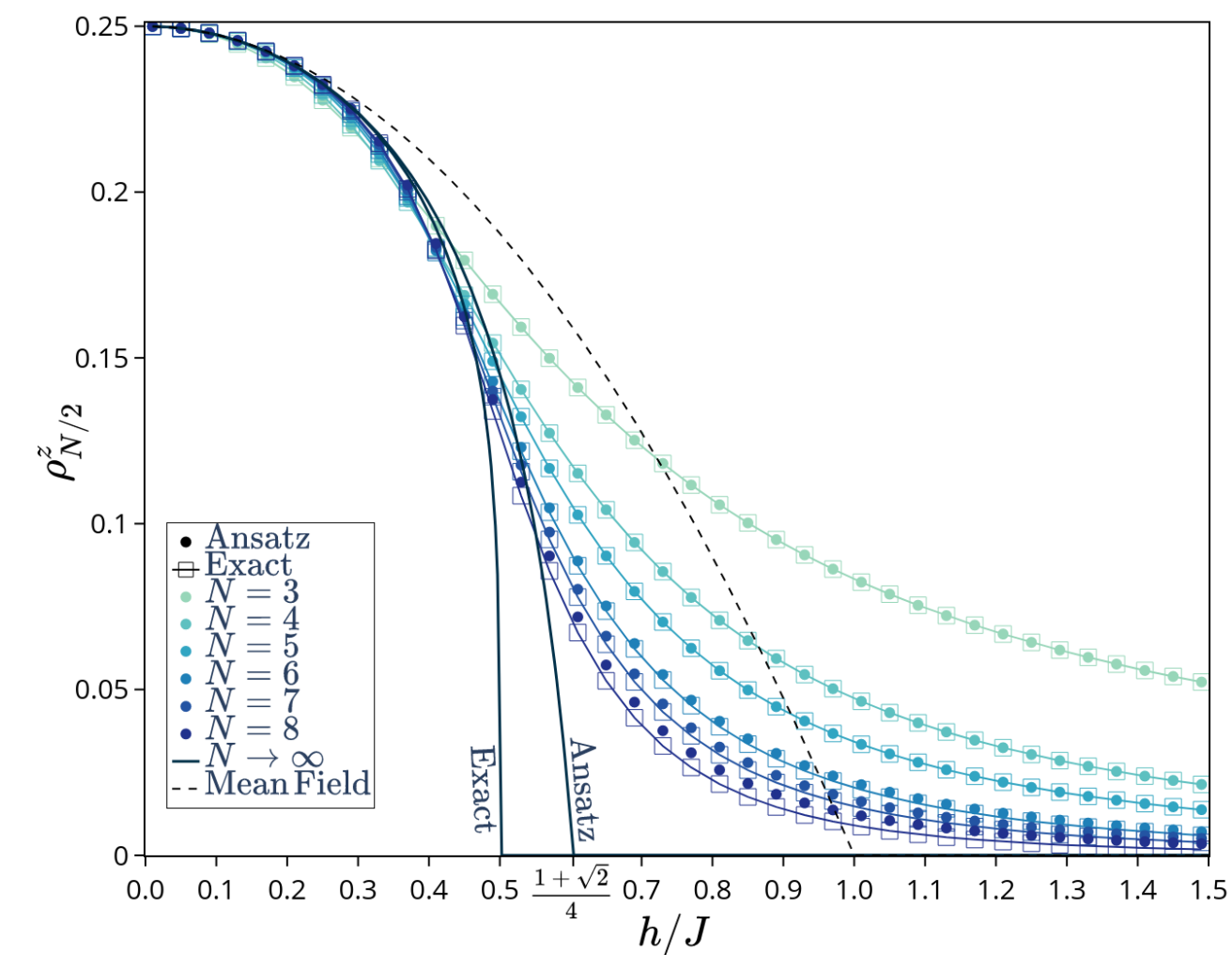


FIG. 5. TFIM long-range correlation  $\rho_{N/2}^z$  vs  $h/J$ . Shows results from the symmetrised ansatz (finite  $N$ ,  $N \rightarrow \infty$ ), exact values, and the mean-field prediction.





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Theoretical and Computational Sciences

# Thank you!

Extra questions? - [sinayskiy@ukzn.ac.za](mailto:sinayskiy@ukzn.ac.za)

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